

# Lumpabilities and Noninterference in PEPA

*Alberto Casagrande*

# PEPA - Performance Evaluation Process Algebra

A **stochastic process algebra** for modelling systems [\[H90\]](#)

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# PEPA - Performance Evaluation Process Algebra

A **stochastic process algebra** for modelling systems [H90]

$$P ::= (a, \lambda). P \mid P + P \mid P \triangleright_{\mathbf{A}} \triangleleft P \mid P / \mathbf{A} \mid A$$

- ▶ *Actions:*  $(a, \lambda). Q$
- ▶ *Choises:*  $P + Q$
- ▶ *Cooperations:*  $P \triangleright_{\mathbf{A}} \triangleleft Q$
- ▶ *Hidings:*  $P / \mathbf{A}$
- ▶ *Definitions:*  $A = P$

# From Operational Semantics to CTMC

PEPA processes are Continuous-Time Markov Chains

```
P = (a, 2).Q + (b, 3).R + (c, 7).P;  
W = (b, 5).S + (c, 11).W;
```

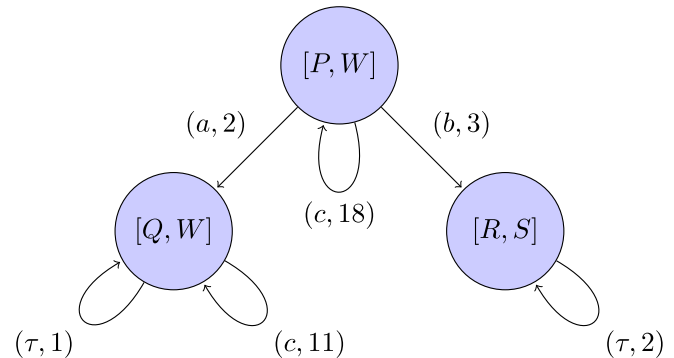
```
Q = (tau, 1).Q;  
S = (tau, 1).S;  
R = (tau, 1).R;
```

```
P<b>W
```

# From Operational Semantics to CTMC

PEPA processes are **Continuous-Time Markov Chains**

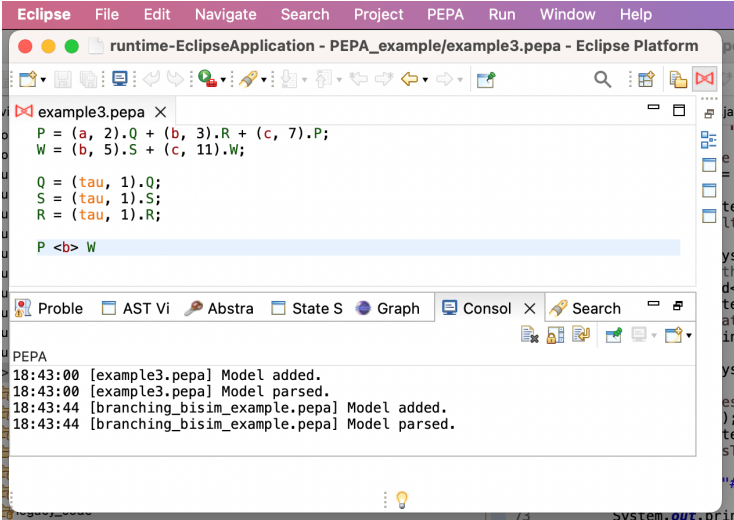
```
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W = (b, 5).S + (c, 11).W;  
  
Q = (tau, 1).Q;  
S = (tau, 1).S;  
R = (tau, 1).R;  
  
P<b>W
```



- nodes are model states
- transitions are (synchronized) actions

# PEPA Implementation

- ▶ model specification
- ▶ steady state analysis
- ▶ scalability analysis



The screenshot shows the Eclipse IDE interface. The top menu bar includes 'Eclipse', 'File', 'Edit', 'Navigate', 'Search', 'Project', 'PEPA', 'Run', 'Window', and 'Help'. The title bar reads 'runtime-EclipseApplication - PEPA\_example/example3.pepa - Eclipse Platform'. The editor window displays the file 'example3.pepa' with the following code:

```
P = (a, 2).Q + (b, 3).R + (c, 7).P;  
W = (b, 5).S + (c, 11).W;  
  
Q = (tau, 1).Q;  
S = (tau, 1).S;  
R = (tau, 1).R;  
  
P <b> W
```

Below the editor, the 'PEPA' console shows the following output:

```
PEPA  
18:43:00 [example3.pepa] Model added.  
18:43:00 [example3.pepa] Model parsed.  
18:43:44 [branching_bisim_example.pepa] Model added.  
18:43:44 [branching_bisim_example.pepa] Model parsed.
```

# **Non-Interference**

Is a property avoiding information leaks due to system behaviours

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- ▶ Non- $\tau$  actions are either **high level** (H) or **low level** (L)
- ▶ **Low level users** can only perform low level actions
- ▶ High actions should be "invisible" to low level users

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What does "invisible" mean?

# Non-Interference

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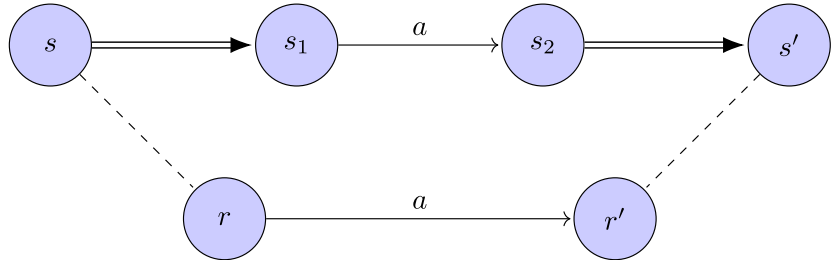
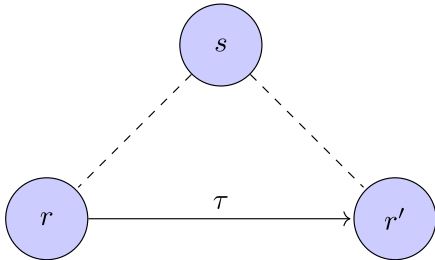
- ▶ Non- $\tau$  actions are either **high level** (H) or **low level** (L)
- ▶ **Low level users** can only perform low level actions
- ▶ High actions should be "invisible" to low level users

What does "invisible" mean? Equivalent w.r.t. some relation

# Weak Bisimulations

$R_w$  is symmetric and if  $(r, s) \in R_w$  and  $r \xrightarrow{a} r'$ , either:

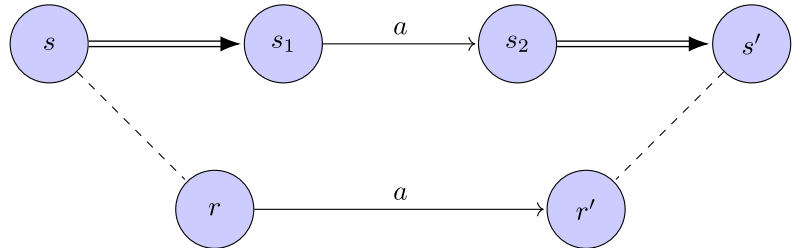
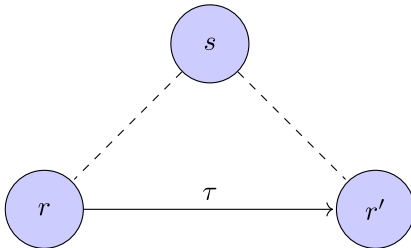
- ▶  $a = \tau$  and  $(r', s) \in R_w$
- ▶  $s \Rightarrow s_1 \xrightarrow{a} s_2 \Rightarrow s'$  and  $(r', s') \in R_w$



# Branching Bisimulations

$R_b$  is symmetric and if  $(r, s) \in R_b$  and  $r \xrightarrow{a} r'$ , either:

- ▶  $a = \tau$  and  $(r', s) \in R_b$
- ▶  $s \Rightarrow s_1 \xrightarrow{a} s_2 \Rightarrow s'$  and  $(r, s_1), (r', s_2), (r', s') \in R_b$

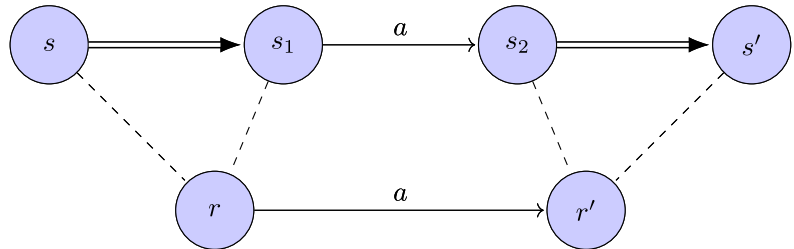
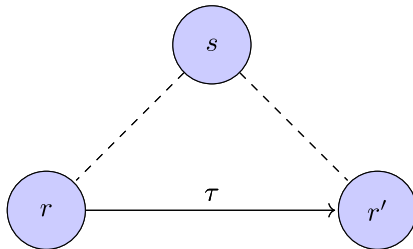


# Branching Bisimulations

$R_b$  is symmetric and if  $(r, s) \in R_b$  and  $r \xrightarrow{a} r'$ , either:

►  $a = \tau$  and  $(r', s) \in R_b$

►  $s \Rightarrow s_1 \xrightarrow{a} s_2 \Rightarrow s'$  and  $(r, s_1), (r', s_2), (r', s') \in R_b$



# Lumpable Bisimulations

[HPR21]  $R_l$  is an equivalence relation on  $\mathcal{P}$  and s.t. if  $(r, s) \in R_l$ , then, for all  $a$  and for all  $C \in \mathcal{P}/R_l$  that either:

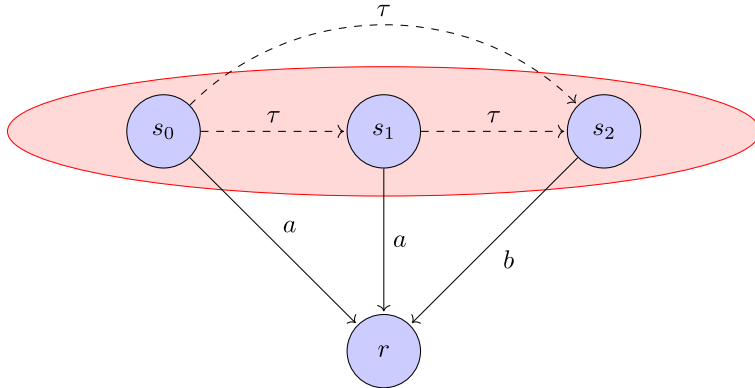
- ▶  $a \neq \tau$  or
- ▶  $a = \tau$  and  $r, s \notin C$ ,

the following equation holds

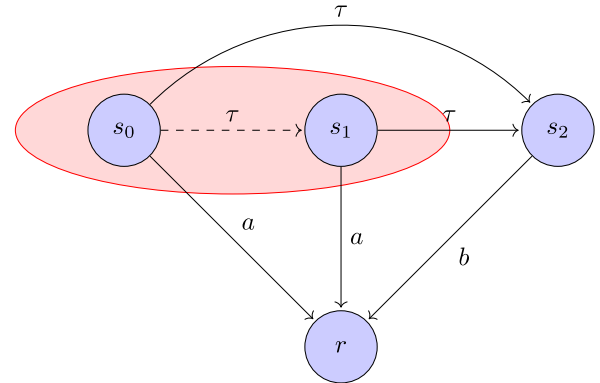
$$\sum_{c \in C} \text{rate} \left( r \xrightarrow{a} c \right) = \sum_{c' \in C} \text{rate} \left( s \xrightarrow{a} c' \right)$$

# Branching vs Lumpable

## Branching



## Lumpable



# Persistent Stochastic Non-Interference

**Definition** [HPR21]

$\forall P'$  s.t.  $P \rightarrow^* P'$  and  $\forall H \in \mathcal{C}_H$

$$P \in \text{PSNI} \quad \Longleftrightarrow \quad (P' \boxtimes_{\mathcal{H}} 0)/\mathcal{H} \approx_l (P' \boxtimes_{\mathcal{H}} H)/\mathcal{H}$$

# Persistent Stochastic Non-Interference

## Definition [HPR21]

$$\begin{array}{l}
 \forall P' \text{ s.t. } P \rightarrow^* P' \text{ and } \forall H \in C_H \\
 P \in \text{PSNI} \quad \Longleftrightarrow \quad (P' \triangleright_{\text{H}} \triangleleft 0)/H \approx_l (P' \triangleright_{\text{H}} \triangleleft H)/H
 \end{array}$$

## Theorem [HPR21]

$$P \in \text{PSNI} \quad \Longleftrightarrow \quad (P \triangleright_{\text{H}} \triangleleft 0)/H \approx_l^H P$$

# Adding PSNI in PEPA Plugin

Two steps

- ▶ action level labelling
- ▶ PSNI testing command

# Adding PSNI in PEPA Plugin

Two steps

- ▶ action level labelling (PEPA language)
- ▶ PSNI testing command (PEPA-plugin UI)

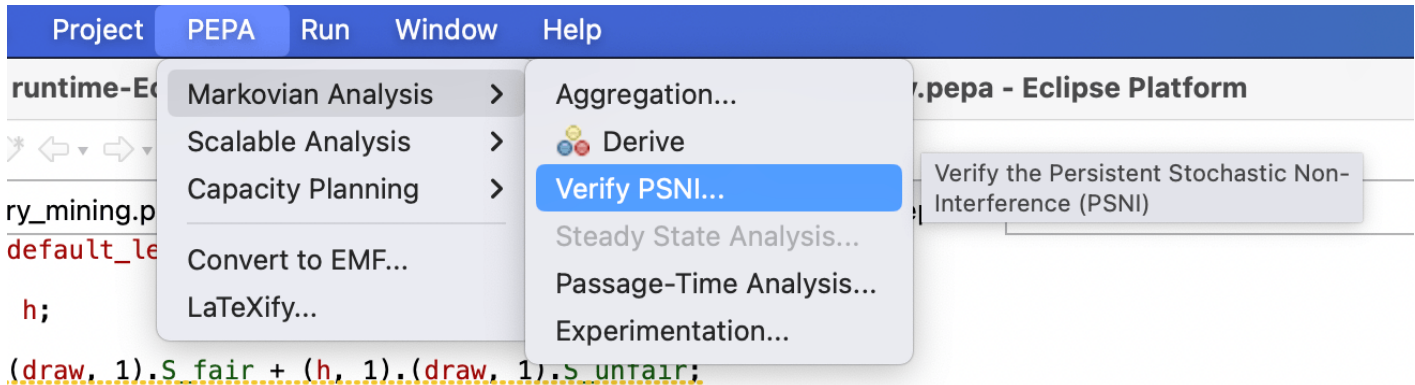
# Action Level Labelling

```
1 ar = 2;  
2 br = 3;  
3  
4 P = (a, ar).Q + (b, br).R + (c, 7).P;  
5 W = (b, 5).S + (c, 11).W;  
6  
7 Q = (tau, 1).Q;  
8 S = (tau, 1).S;  
9 R = (tau, 1).R;  
10  
11 P<b>W
```

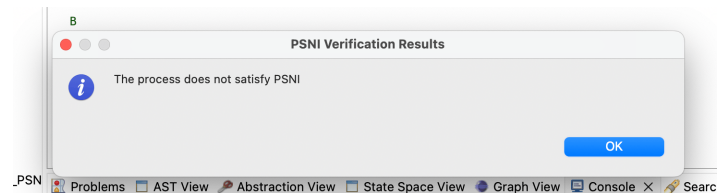
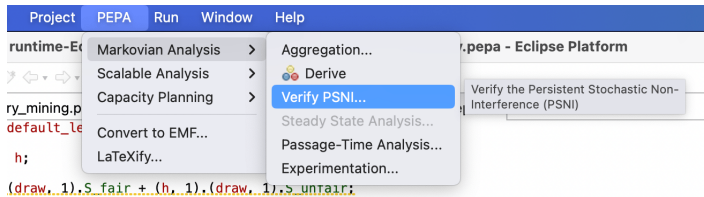
# Action Level Labelling

```
1 ar = 2;  
2 br = 3;  
3  
4 set_default_level low;  
5  
6 high a, b;  
7  
8 P = (a, ar).Q + (b, br).R + (c, 7).P;  
9 W = (b, 5).S + (c, 11).W;  
10  
11 Q = (tau, 1).Q;  
12 S = (tau, 1).S;  
13 R = (tau, 1).R;  
14  
15 P<b>W
```

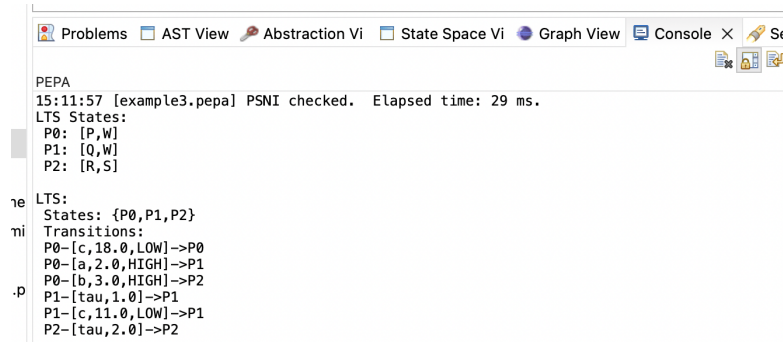
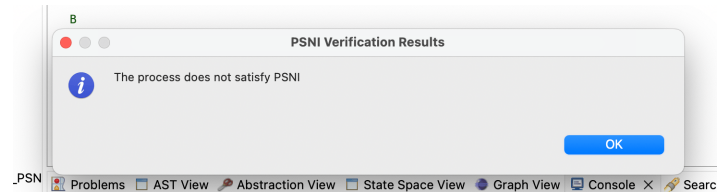
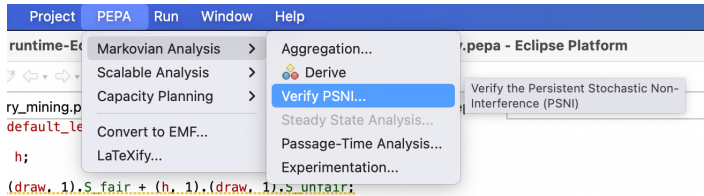
# PEPA Plug-in Integration



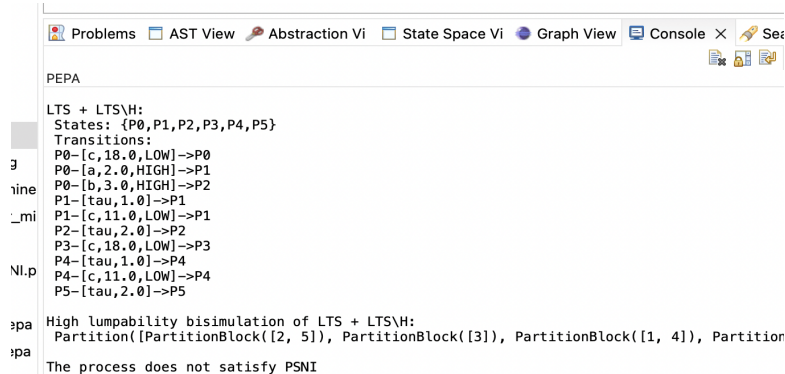
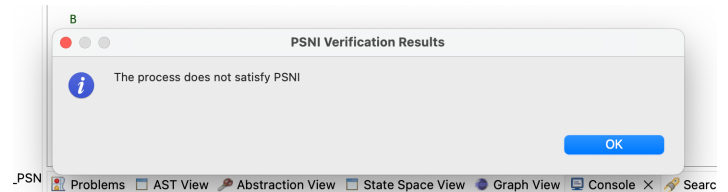
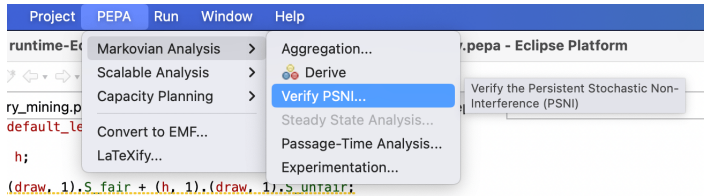
# PEPA Plug-in Integration



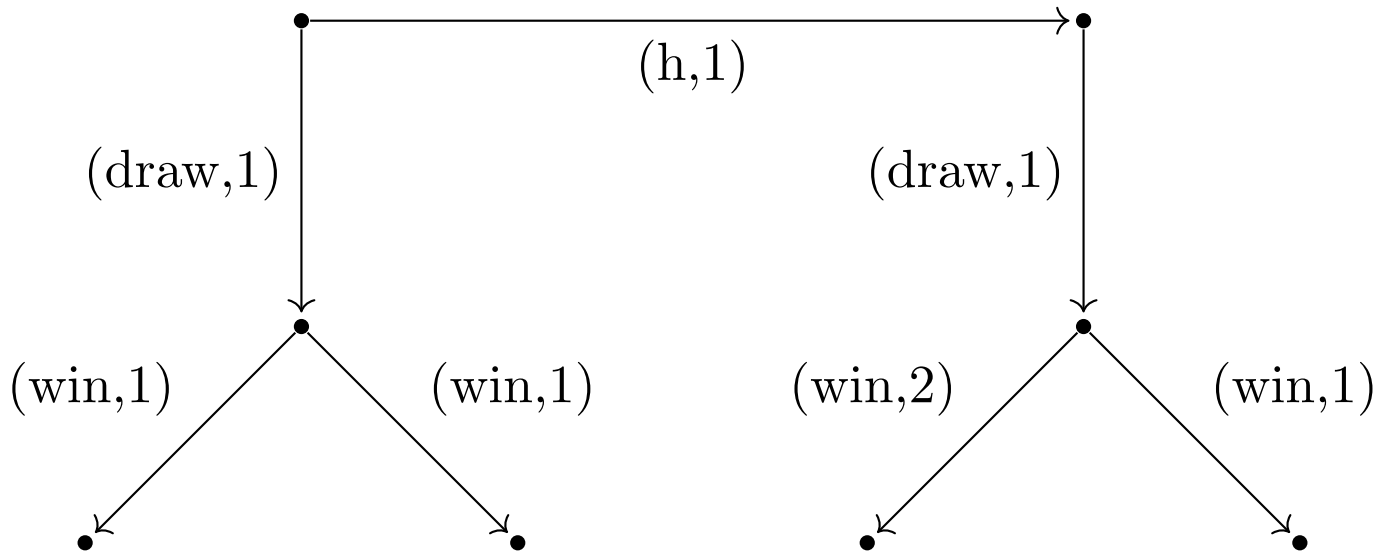
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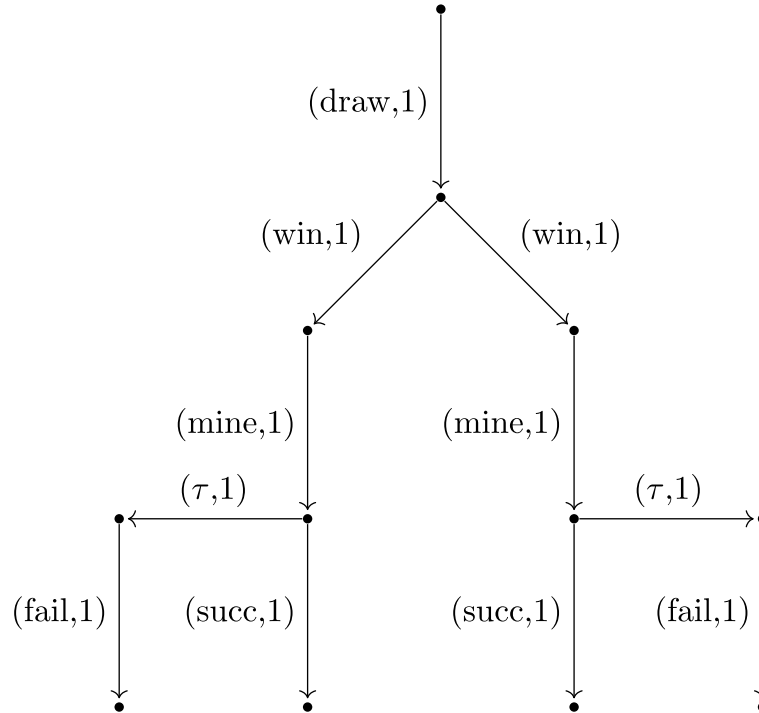
# Example I: Non-Uniform Lottery



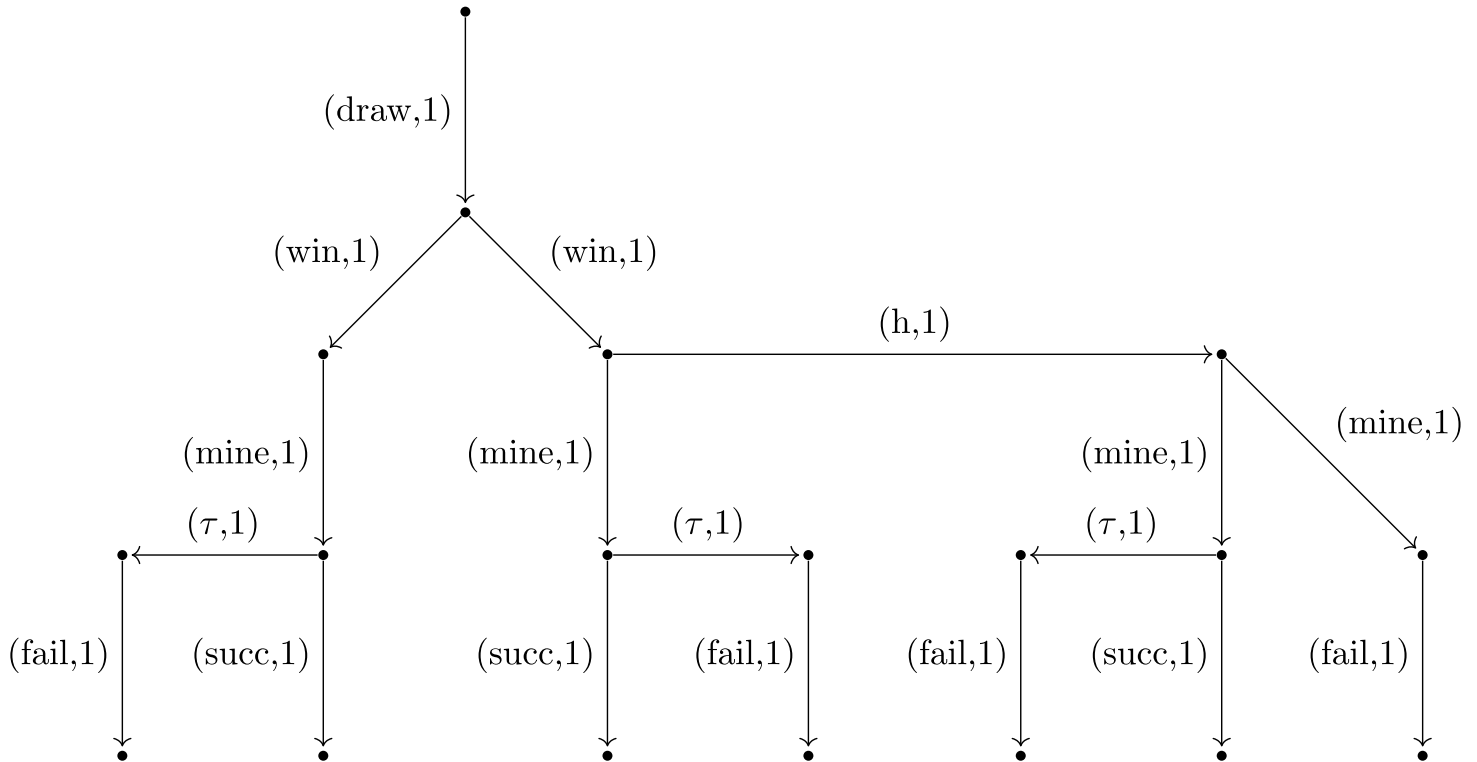
# Example I: Non-Uniform Lottery

```
1 set_default_level low;
2
3 high h;
4
5 B0 = (draw, 1).D0 + (h, 1).(draw, 1).D1;
6
7 D0 = (win, 1).W0 + (win, 1).W1;
8
9 W0 = (tau, 1).W0;
10 W1 = (tau, 1).W1;
11
12 D1 = (win, 2).W2 + (win, 1).W3;
13
14 W2 = (tau, 1).W2;
15 W3 = (tau, 1).W3;
16
17 B0
```

## Example 2: Fair Lottery Verifier



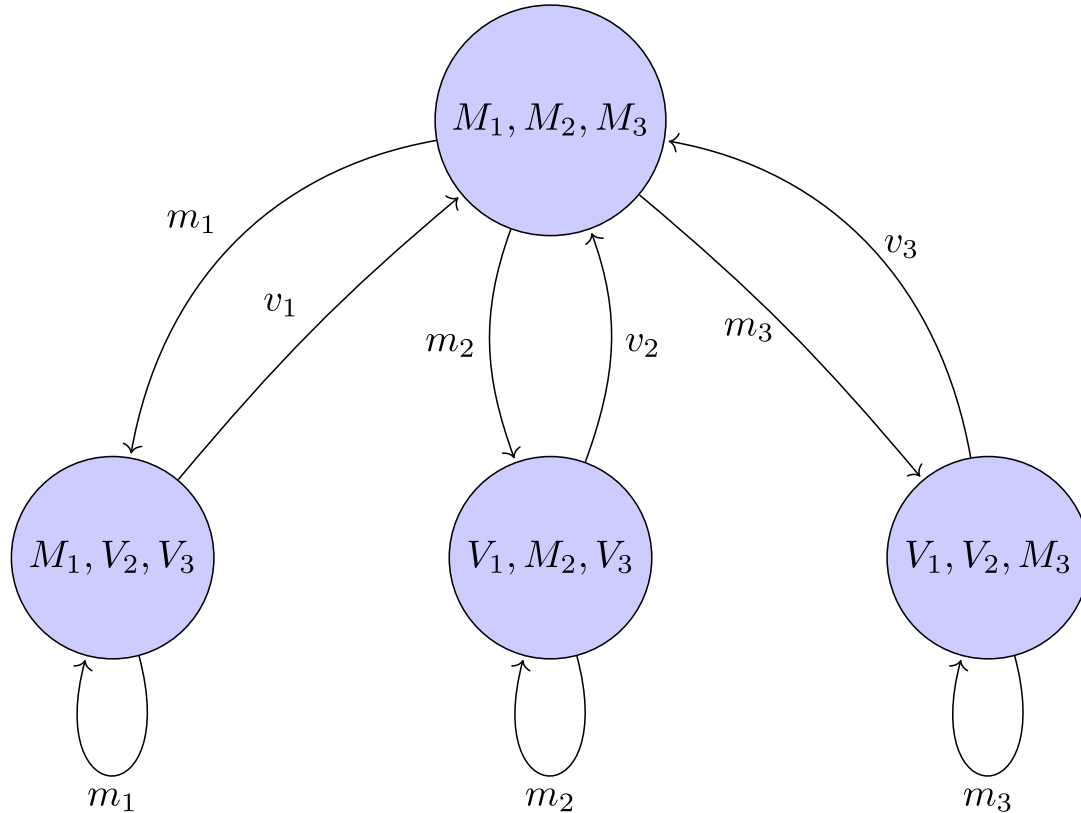
## Example 2: Unfair Lottery Verifier



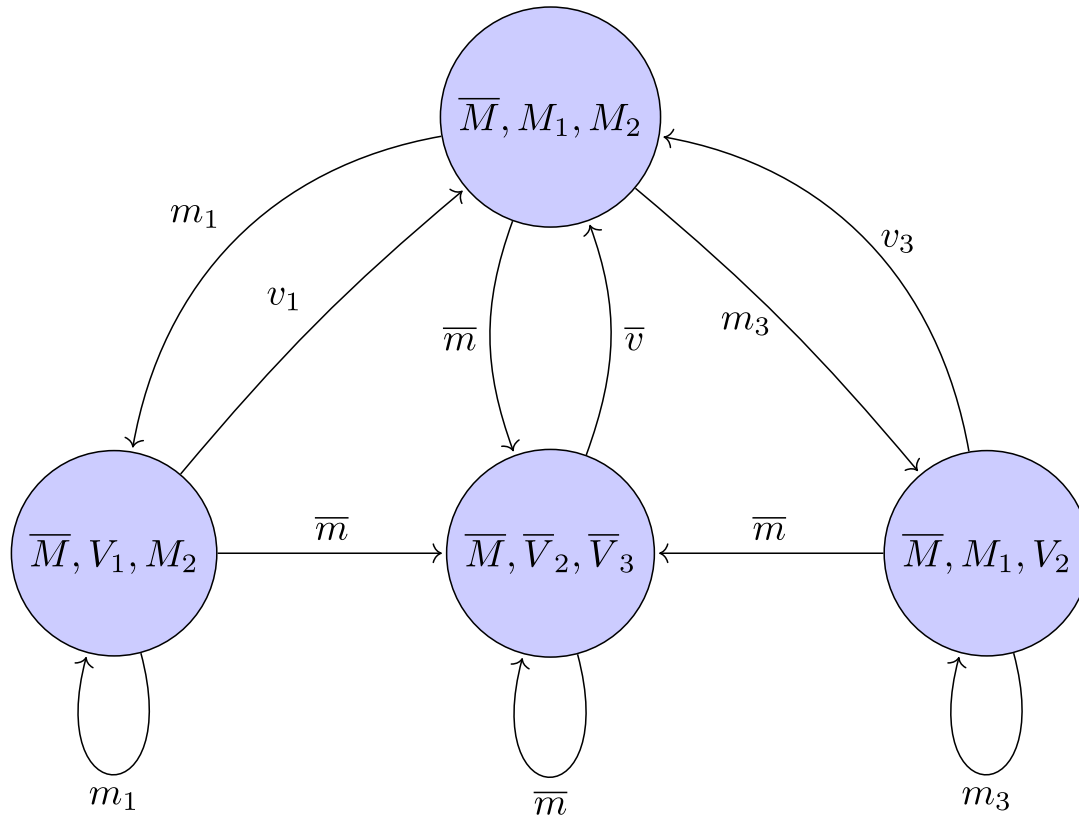
## Example 2: Unfair Lottery Verifier

```
1 set_default_level low;
2 high h;
3
4 B = (win0, 1).W0 + (win1, 1).W1;
5
6 W0 = (mine, 1).((tau,1).(fail,1).F0 + S0);
7
8 W1 = (mine, 1).((tau,1).(fail,1).F1 + S1)
9       + (h, 1).((mine,1).(fail, 1).F2
10                + (mine, 1).((tau,1).(fail,1).F3 + S3));
11
12 F0 = (tau, 1).F0;
13 S0 = (tau, 1).S0;
14 ...
15 F3 = (tau, 1).F3;
16 S3 = (tau, 1).S3;
17
18 B
```

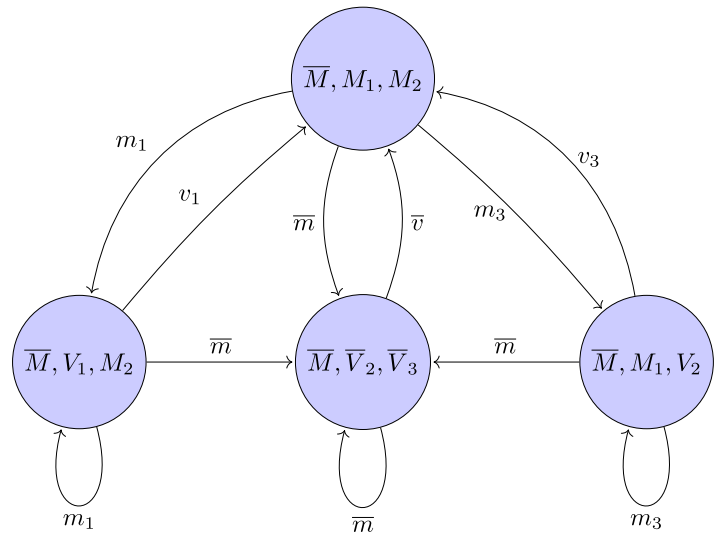
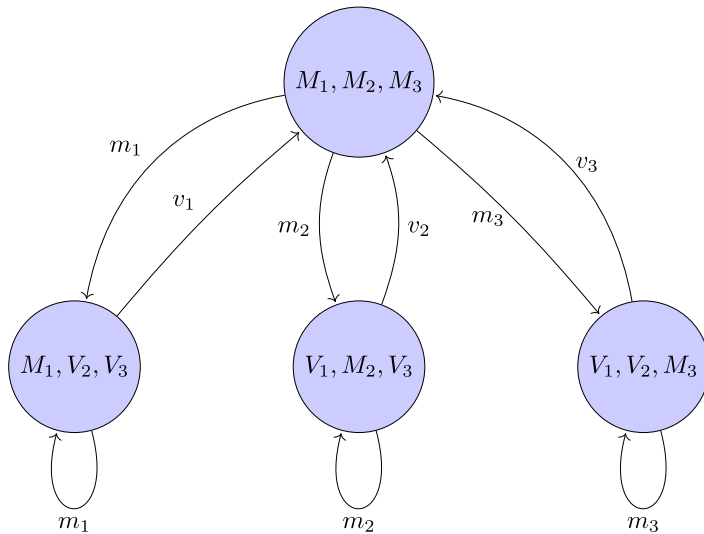
## Example 3: Fair Miners [SMMvMR22]



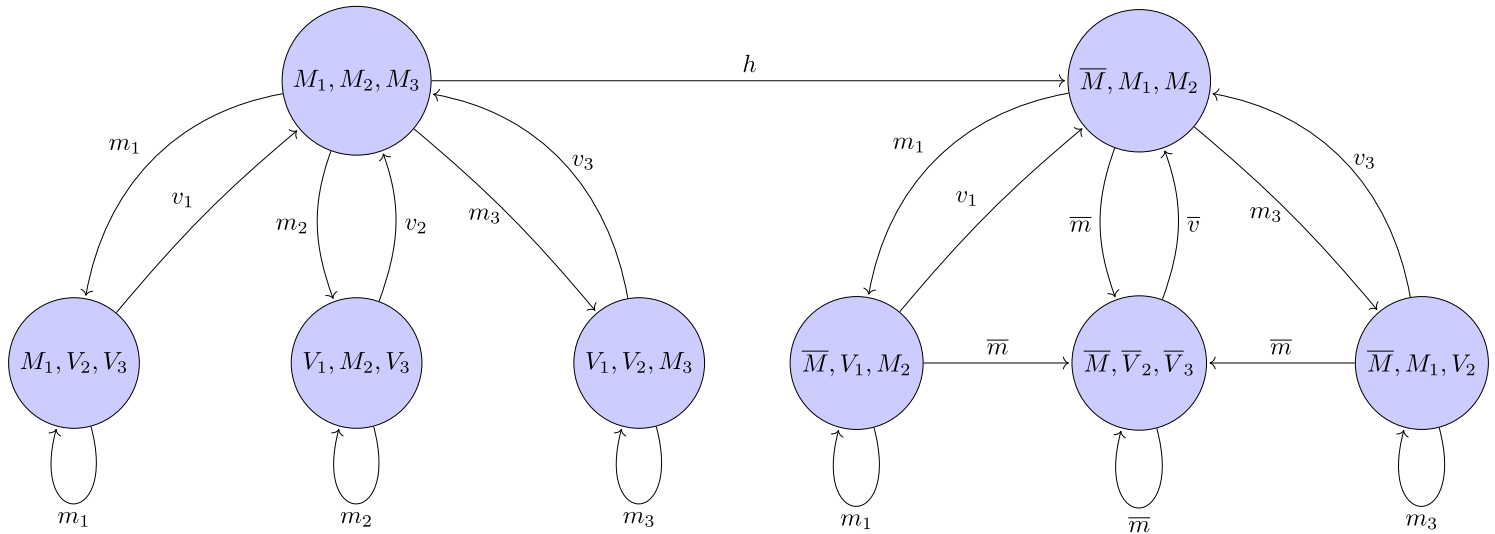
# Example 3: Unfair Miner [SMMvMR22]



# Example 3: Fair-Unfair Miner [SMMvMR22]



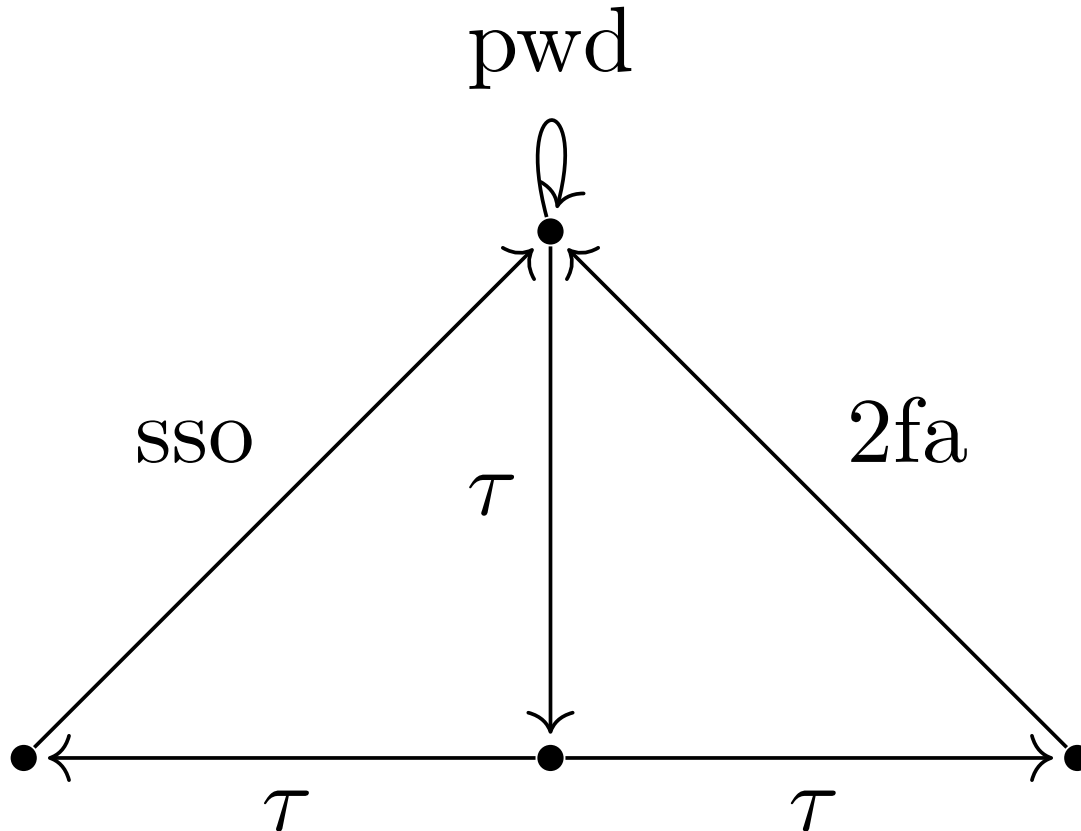
# Example 3: Fair-Unfair Miner [SMMvMR22]



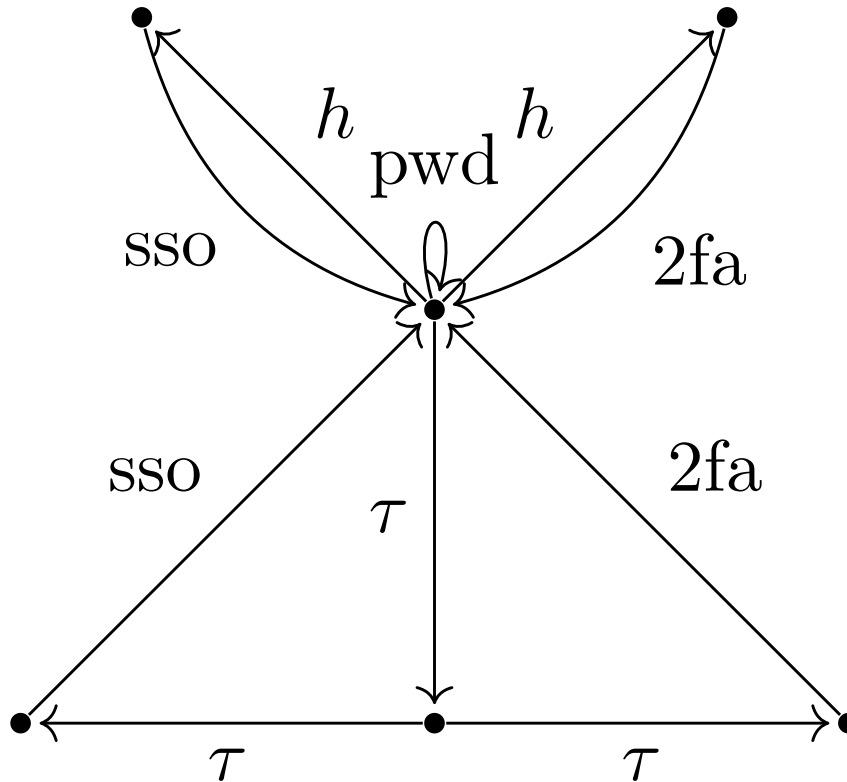
## Example 3: Fair-Unfair Miner [SMMvMR22]

```
1 set_default_level low;
2
3 high h;
4
5 // 3 fair miners + h transition
6 P0 = (m1,0.5).P1 + (m2,0.5).P2 + (m3,0.5).P3 + (h, 0.1).P4;
7 P1 = (v1,0.3).P0 + (m1,0.5).P1;
8 P2 = (v2,0.3).P0 + (m2,0.5).P2;
9 P3 = (v3,0.3).P0 + (m3,0.5).P3;
10
11 // 2 fair miners, 1 unfair
12 P4 = (mu,0.5).P5 + (m1,0.5).P6 + (m2,0.5).P7;
13 P5 = (vu,0.3).P4 + (mu,0.5).P5;
14 P6 = (v1,0.3).P4 + (mu,0.5).P5 + (m1,0.5).P6;
15 P7 = (v2,0.3).P4 + (mu,0.5).P5 + (m2,0.5).P7;
16
17 P0
```

## Example 4: DBMS Auth [EABR23]



## Example 4: DBMS Auth Attack [EABR23]



## Example 4: DBMS Auth Attack [\[EABR23\]](#)

```
1 set_default_level low;
2
3 high h;
4
5 WT0 = (l_pwd, 1).WT0
6       + (tau, 1).((tau, 1).(l_sso, 1).WT0
7       + (tau, 1).(l_2fa, 1).WT0)
8       + (h, 1).(l_sso, 1).WT0
9       + (h, 1).(l_2fa, 1).WT0;
10
11 WT0
```

**Thank you!**