Enhancing PEPA Eclipse Plugin

Riccardo Romanello¹

¹ Ca' Foscari, University of Venice

May 30, 2025









Background

▶ PEPA is Performance Evaluation Process Algebra







Background

- ▶ PEPA is Performance Evaluation Process Algebra
- Such theoretical framework has been deeply investigated in the literature







Background

- ▶ PEPA is Performance Evaluation Process Algebra
- Such theoretical framework has been deeply investigated in the literature
- ► PEPA Eclipse Plugin serves as a tool for testing and visualizing the results of such investigation







▶ We increased the set of features of the PEPA Eclipse Plugin







- ▶ We increased the set of features of the PEPA Eclipse Plugin
- ► We added a (time) reversibility control







- ▶ We increased the set of features of the PEPA Eclipse Plugin
- We added a (time) reversibility control
- We implemented all the mechanisms to handle PSNI







- ▶ We increased the set of features of the PEPA Eclipse Plugin
- We added a (time) reversibility control
- We implemented all the mechanisms to handle PSNI
- ► To conclude, we introduced the notion of reversible actions







▶ What I will not explain









- ► What I will not explain
- ► PEPA 101







- What I will not explain
- ► PEPA 101
- PEPA Eclipse Plugin
 - Reversible processes







- What I will not explain
- ► PEPA 101
- PEPA Eclipse Plugin
 - Reversible processes
 - Reversible actions







- What I will not explain
- ► PEPA 101
- PEPA Eclipse Plugin
 - Reversible processes
 - Reversible actions
 - High and Low level actions







Assumptions

I assume you all have heard, at least once, about the following topics:

- Stochastic Processes
- Markov Chains (either Discrete or Continuous Time)
- Markov Property
- Representation of Markov Chains









Encoding stochastic processes

- ► The description of stochastic processes via Markov chains is a cumbersome and often unfeasible activity
- Moreover, it is hard to check if the stochastic process is correct or not
- Adding an abstraction layer by using a compact representation that is easy to produce and verify, can be a way to address both the issues
- Stochastic process algebras fit this role of additional layer
- We adopt PEPA as process algebra to model stochastic processes.









Performance Evaluation Process Algebra

- ► A PEPA model consists of a set of components that engage either individually or cooperatively in activities
- Components represent identifiable units, and can be either atomic or composed
- Activities capture actions performed by the components.
 Every activity is associated with an action type
- Activities are not instantaneous. The probability that an activity a with rate r happens within a period of time of length t is given by $F_a(t) = 1 e^{-rt}$
- ► Hence, an activity with action type α and rate r is completely defined by the pair (α, r) .







PEPA - Syntax and Semantics

Definition - PEPA Syntax

Let \mathcal{A} be a set of actions with $\tau \in \mathcal{A}$ Let $\alpha \in \mathcal{A}$, $A \subseteq \mathcal{A}$, and $r \in \mathbb{R} \cup \{\top\}$

$$S ::= \mathbf{0} \mid (\alpha, r).S \mid S + S \mid X$$

$$P ::= P \bowtie_{A} P \mid P/A \mid P \setminus A \mid S$$

Each variable X is associated to a definition $X \stackrel{\text{def}}{=} P$

PEPA defines Labeled Continuous Time Markov Chains

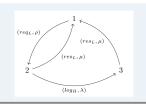








Example



$$X_1 = (req_L, \rho).X_2$$

 $X_2 = (res_L, \mu).X_1 + (log_H, \lambda).X_3$
 $X_3 = (res_L, \mu).X_1$







PEPA - Semantics for Synchronization

$$\frac{P \xrightarrow{(\alpha,r)} P'}{P \bowtie_{A} Q \xrightarrow{(\alpha,r)} P' \bowtie_{A} Q} (\alpha \notin A) \qquad \frac{Q \xrightarrow{(\alpha,r)} Q'}{P \bowtie_{A} Q \xrightarrow{(\alpha,r)} P \bowtie_{A} Q'} (\alpha \notin A)$$

$$\frac{P \xrightarrow{(\alpha,r_{1})} P' Q \xrightarrow{(\alpha,r_{2})} Q'}{P \bowtie_{A} Q \xrightarrow{(\alpha,r_{2})} Q'} \quad (\alpha \in A)$$

$$\text{where } R = \frac{r_{1}}{r_{\alpha}(P)} \frac{r_{2}}{r_{\alpha}(Q)} \min(r_{\alpha}(P), r_{\alpha}(Q))$$







- Reversibility of Markovian Processes can be modelled in two ways
- ► Time reversibility deals with inverting the flow of time
- Causal reversibility deals with undoing operations in the correct order







- ► Time Reversibility is usually characterized by means of two results:
 - Kolmogorov condition
 - Detailed Balance Equations







ρ-Reversibility [3]

```
Algorithm 1: ReversibleUpTo(S, P, \rho)

for i \in S do
|color[i] = white
|I[i] = +\infty
end
I[1] = 1
return DFS-ReversibleUpTo(S, P, 1, \rho)
```

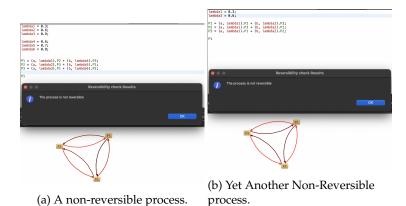
```
Algorithm 2: DFS-ReversibleUpTo(S, P, i, \rho)
 bool = true
 for bool \wedge p[i,j] \neq 0 do
     if (p[\rho[j], \rho[i]] = 0) then
      | bool = false
     end
     if (color[j] \neq white \land \Pi[i]p[i,j] \neq \Pi[j]p[\rho[j], \rho[i]])
      then
      | bool = false
      end
     if (color[i] = white) then
         color[j] = grey
         \Pi[j] = \Pi[i] \frac{p[i,j]}{p[\rho[j],\rho[i]]}
         bool = bool \land DFS-ReversibleUpTo(\mathcal{S}, \mathbf{P}, i, \rho)
     end
 end
 return bool
```







What we did











A Positive Case

```
lambda1 = 0.3;
 lambda2 = 0.4;
 lambda3 = 0.5;
P1 = (a, lambda1).P2 + (b, lambda3).P3;
P2 = (a, lambda2).P3 + (b, lambda1).P1;
P3 = (a, lambda3).P1 + (b, lambda2).P2;
                                                      Reversibility check Results
            The process is reversible
```









Another Positive Case

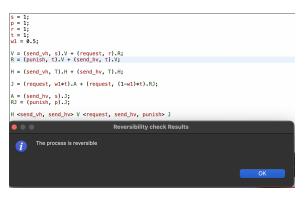


Figure 4: An (Hacked) Blockchain with punish mechanism. Process defined in a yet to be published work.







▶ The second feature deals with *reversibility*







- ▶ The second feature deals with *reversibility*
- Starting from the previous work presented in [1]







- The second feature deals with reversibility
- Starting from the previous work presented in [1]
- ► We defined a set of reversible actions







- The second feature deals with reversibility
- Starting from the previous work presented in [1]
- We defined a set of reversible actions
- ► If a process is reached via a reversible action, it cannot be reached by any other action







- The second feature deals with reversibility
- ▶ Starting from the previous work presented in [1]
- ► We defined a set of reversible actions
- ► If a process is reached via a reversible action, it cannot be reached by any other action
- ► In the LTS we introduce a *backward* edge per each reversible transition







A Legal Reversible Action

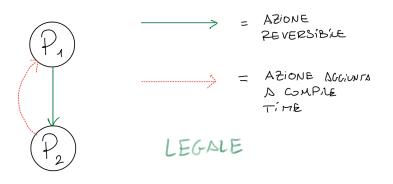


Figure 5: A correctly used reversible action









An Illegal Reversible Action

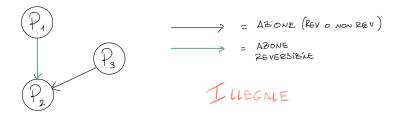


Figure 6: An incorrectly used reversible action



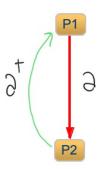




A simple reversible tree

```
one = 1.0;
rev_actions a;
P1 = (a, one).P2;
P1
```

(a) PEPA process.



(b) (Edited) induced LTS.





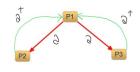




Yet Another Simple Reversible Tree

```
one = 1.0;
rev_actions a;
P1 = (a,_one),P2 + (a,_one),P3;
P1
```

(a) PEPA process.



(b) (Edited) induced LTS.







► The last feature concerns Persistent Stochastic Non-Interference







- ► The last feature concerns Persistent Stochastic Non-Interference
- ► Non-Interference aims at protecting sensitive data from undesired accesses







- ► The last feature concerns Persistent Stochastic Non-Interference
- Non-Interference aims at protecting sensitive data from undesired accesses
- ► Goguen-Meseguer'82: information does not flow from high (confindential) to low (public) if the high behavior cannot be observed at low level







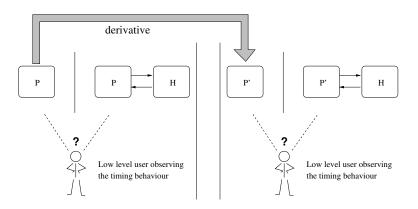
- ► The last feature concerns Persistent Stochastic Non-Interference
- Non-Interference aims at protecting sensitive data from undesired accesses
- Goguen-Meseguer'82: information does not flow from high (confindential) to low (public) if the high behavior cannot be observed at low level
- Persistency: Non-Interference has to be guaranteed in all the states of the system, if processes migrate during execution







Intuitively



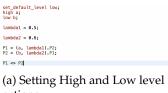




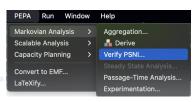




An example



actions



(b) Verifying PSNI







A Lottery Ticket Example

```
one = 1.0;
half = 0.5:
epsilon = 0.000001;
set_default_level low;
high drawh;
F = (w1, half).W1 + (w2, half).W2;
UF = (tau, 1).UF1;
UF1 = (w1, one-epsilon).W1 + (w2, epsilon).W2;
W1 = (tau, one).W1;
W2 = (tau, one).W2;
P = (tau, one).F + (drawh, one).UF + (tau, one).P;
lΡ
```

Figure 10: A Non-PSNI process, modelled from an example provided in [2].









- [1] Marco Bernardo and Sabina Rossi.
 Reverse bisimilarity vs. forward bisimilarity.
 In Orna Kupferman and Pawel Sobocinski, editors,
 Foundations of Software Science and Computation
 Structures, pages 265–284, Cham, 2023. Springer Nature
 Switzerland.
- [2] Andrea Esposito, Alessandro Aldini, and Marco Bernardo. Noninterference analysis of irreversible or reversible systems with nondeterminism and probabilities, 2025.
- [3] Matteo Sottana, Carla Piazza, and Andrea Albarelli. Efficient computation of renaming functions for p-reversible discrete and continuous time markov chains. In Proceedings of the 11th EAI International Conference on Performance Evaluation Methodologies and Tools,







VALUETOOLS 2017, page 52–59. Association for Computing Machinery, 2017.







