# **Exact Non-Interference**

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- Current approaches [2] are built upon the notion of Lumpable Bisimulation, aka Strong Equivalence
- ▶ We want to obtain similar results using *Exact Equivalence*
- Roughly speaking:
  - Strong: looks at ougoing rates
  - Exact: looks at incoming rates







#### Outline of the Talk

- ► An High-Level view of PSNI
- Observational Equivalence
- PSNI and strong equivalence
- Exact Equivalence
- Weak-exact equivalence







#### Persistent Stochastic Non-Interference

#### The Context

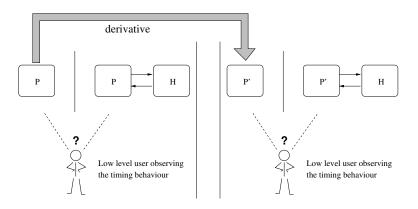
- ► Non-Interference aims at protecting sensitive data from undesired accesses
- Goguen-Meseguer'82: information does not flow from high (confindential) to low (public) if the high behavior cannot be observed at low level
- Persistency: Non-Interference has to be guaranteed in all the states of the system, if processes migrate during execution







### Intuitively



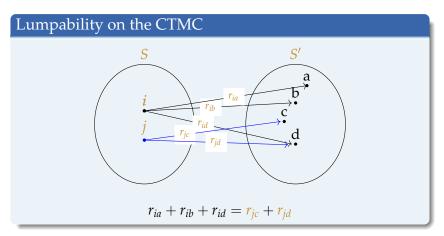








# Observation Equivalence



Users cannot distinguish lumpable bisimilar PEPA components









# Observation Equivalence

### Definition - Lumpable bisimilarity

It is the largest equivalence relation  $\approx_l$  such that if  $P \approx_l Q$ , then for all  $\alpha$  and for each S equivalence class

- either  $\alpha \neq \tau$ ,

it holds

$$\sum_{P' \in S, \ P \xrightarrow{(\alpha, r_{\alpha})} P'} r_{\alpha} = \sum_{Q' \in S, \ Q \xrightarrow{(\alpha, r_{\alpha})} Q'} r_{\alpha}$$

It is contextual, action preserving, and induces a lumpability







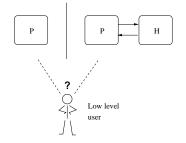


#### Non-Interference

### A general definition [Focardi-Gorrieri'95]

 $P \in NI$  iff  $\forall$  high level process H,  $(P|0) \sim^{low} (P|H)$ 

where  $\sim^{low}$  denotes a low level observation equivalence











### Stochastic Non-Interference (SNI)

- ▶ We partition the actions into  $\mathcal{L}$  (low),  $\mathcal{H}$  (high),  $\{\tau\}$  (synch.)
- High (low) level processes can only perform(/observe) high (low) level actions

#### **Definition - SNI**

 $P \in SNI$  iff  $\forall$  high level PEPA component H

$$(P \bowtie_{\mathcal{H}} 0) \sim^{low} (P \bowtie_{\mathcal{H}} H)$$

The above can be rewritten as:

$$(P \bowtie_{\mathcal{H}} 0)/\mathcal{H} \approx_{l} (P \bowtie_{\mathcal{H}} H)/\mathcal{H}$$







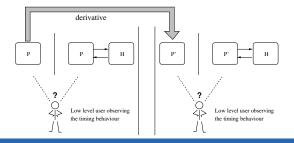


### Persistent SNI (PSNI)

#### **Definition - PSNI**

 $P \in PSNI$  iff  $\forall$  derivative P' of P

 $P' \in SNI$ 



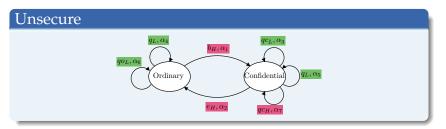


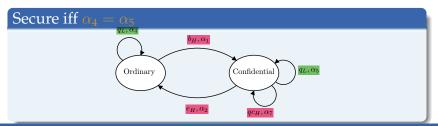






# Toy Example: Unsecure Vs Secure System













### Our new approach

- ► The above notion of PSNI is based on strong equivalence/lumpable bisimilarity
- ▶ What if we leverage such notion to exact equivalence?
- We are defining a weak exact equivalence that treats  $\tau$  actions in different ways that classical exact equivalence does
- Addressing the problems of:
  - Formalize PSNI in terms of such new notion
  - Provide an efficient algorithm to test this newly introduced concept







# Exact Equivalence

### Exact Equivalence on CTMC

Let X(t) be a CTMC with state space  $S = \{0, 1, ..., n\}$  and  $\sim$  be an equivalence relation over S. We say that X(t) is exactly lumpable with respect to  $\sim$  if for any [k],  $[l] \in S / \sim$  and  $i, j \in [l]$ , it holds that  $q_{[k],i} = q_{[k],j}$ .

### Exact Equivalence on PEPA components

An equivalence relation over PEPA components,  $\mathcal{R} \subseteq \mathcal{C} \times \mathcal{C}$  is an *exact equivalence* if whenever  $(P,Q) \in \mathcal{R}$ , then for all  $\alpha \in \mathcal{A}$ :

- ▶  $q[S, P, \alpha] = q[S, Q, \alpha] \quad \forall S \in C/\mathcal{R}$









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- The one on PEPA components seems stricter than the other on CTMCs
- ▶ The condition on total outgoing rates was introduced in [1]
- ► It must be introduced because of the diagonal elements in the infinitesimal generator
- That have no counterpart in the PEPA settings
- ► The outgoing rate condition is fundamental for the equiprobability of the partition induced by the relation







### Exact Equivalence, **7** transitions, and PSNI

- ► Recalling that the property we want to ensure is the following:  $(P \bowtie_{\mathcal{H}} 0)/\mathcal{H} \approx_l (P \bowtie_{\mathcal{H}} H)/\mathcal{H}$
- ▶ It is fundamental to carefully deal with  $\tau$ -transitions







### Exact Equivalence, $\tau$ transitions, and PSNI

- ► Recalling that the property we want to ensure is the following:  $(P \bowtie_{\mathcal{H}} 0)/\mathcal{H} \approx_l (P \bowtie_{\mathcal{H}} H)/\mathcal{H}$
- ▶ It is fundamental to carefully deal with  $\tau$ -transitions
- Lumpable bisimulation achieves such endeavor by imposing different conditions according to the symbol  $\alpha$
- Defining PSNI with exact equivalence would have led to a trivial property







# Weak Exact Equivalence

### Weak Exact Equivalence

An equivalence relation over PEPA components,  $\mathcal{R} \subseteq \mathcal{C} \times \mathcal{C}$  is a *weak exact equivalence* if whenever  $(P,Q) \in \mathcal{R}$ , then for all  $\alpha \in \mathcal{A}$ :

- either  $\alpha \neq \tau$  and;
  - $q[P, \alpha] = q[Q, \alpha]$
  - $q[S, P, \alpha] = q[S, Q, \alpha] \quad \forall S \in \mathcal{C}/\mathcal{R}$
- ightharpoonup or  $\alpha = \tau$  and:
  - $q[S, P, \alpha] = q[S, Q, \alpha] \quad \forall S \in \mathcal{C}/\mathcal{R} \cdot P, Q \notin S$







#### The results so far

- We were able to prove that our new notion induces an equiprobable stationary distribution in the underlying CTMC
- ► This implies that the reversed CTMC has a strong equivalence
- ► We are now mimicking the proofs for PSNI by substituting strong equivalence with the weak exact equivalence.







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