

***THE SLOW-SCALE STOCHASTIC
SIMULATION ALGORITHM***

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DG Current Support: **U.C. Santa Barbara
Caltech (NIGMS , NIH)**

DG Past Support: **Caltech (DARPA/AFOSR, Beckman)
Molecular Sciences Institute (Sandia / DOE)
ONR**

ASSUMPTIONS & DEFINITIONS

- A **well-stirred** chemical system at constant volume and temperature.
- N **species** $\{S_1, \dots, S_N\}$. System **state** is $\mathbf{X}(t) = (X_1(t), \dots, X_N(t))$,
 $X_i(t) \equiv$ number of S_i molecules at time t .
- M **reactions** $\{R_1, \dots, R_M\}$. Each R_j is described by two quantities:
 - **State change vector:** $\boldsymbol{\nu}_j = (\nu_{1j}, \dots, \nu_{Nj})$, where
 $\nu_{ij} \equiv$ change induced in X_i by one R_j event.
So R_j induces $\mathbf{x} \rightarrow \mathbf{x} + \boldsymbol{\nu}_j$.
 - **Propensity function:** $a_j(\mathbf{x})$, where
 $a_j(\mathbf{x})dt \equiv$ probability, given $\mathbf{X}(t) = \mathbf{x}$, that R_j will fire in $[t, t + dt)$.

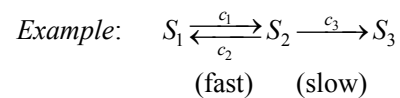
The Stochastic Simulation Algorithm (SSA)

- If the system is in state \mathbf{x} at time t , then with $a_0(\mathbf{x}) \equiv \sum_{k=1}^M a_k(\mathbf{x})$,
 - the time τ to the next reaction will be the exponential random variable with mean $1/a_0(\mathbf{x})$;
 - the index j of the next reaction will be the integer random variable with probability $a_j(\mathbf{x})/a_0(\mathbf{x})$.
- By generating such random numbers for τ and j , we can advance the system to the next reaction:

$$t \leftarrow t + \tau \text{ and } \mathbf{x} \leftarrow \mathbf{x} + \boldsymbol{\nu}_j.$$

A “Multi-scale” Problem

- If the various reaction channels fire at *vastly different rates*, the SSA will spend most of its time simulating the fastest reactions.



- If, in addition, the fastest reactions are “less important” than the slower ones, the SSA simulation will be slow and inefficient.
 - Is there an approximate but accurate way to *skip over* the fast reactions and simulate *only the slow* ones?
 - And how can we tell when it’s okay to do that?

- STIFF SYSTEMS -

- In ODE theory, a *stiff* system is one that evolves on *wide range of time scales*, with the *fastest* mode being *stable*.
 - Approximate methods have been devised to efficiently simulate stiff ODEs.
- But how can we efficiently simulate *stochastic stiff systems*?

SOME ISSUES

- Do “fast” and “slow” apply to *reactions* or to *species*? Or to both?
- How should fast and slow reactions/species be *defined*?
- Since the fast and slow components of a system are *interconnected*, can they be teased apart without seriously altering their combined effects?

The SLOW-SCALE SSA

- A computational recipe for simulating the stochastic evolution of a stiff system *one slow reaction at a time*.
- Proceeds in a *series of steps*, which also serve to *authenticate* the procedure.

1st Provisionally partition the reactions $\{R_1, \dots, R_M\}$:

- Fast reactions $\{R_1^f, \dots, R_{M_f}^f\}$. Slow reactions $\{R_1^s, \dots, R_{M_s}^s\}$.
- Tentative criterion: Want $a_j^f(\mathbf{x}) \gg a_{j'}^s(\mathbf{x})$ ($\forall j, j'$) “most of the time”.
- *The partitioning will later be subjected to an acceptance test.*
- If *no* partitioning is acceptable, then we’re SOL.

2nd Partition the species $\{S_1, \dots, S_N\}$:

- Fast species $\{S_1^f, \dots, S_{N_f}^f\} \Rightarrow \mathbf{X}^f(t) \equiv (X_1^f(t), \dots, X_{N_f}^f(t))$.
- Slow species $\{S_1^s, \dots, S_{N_s}^s\} \Rightarrow \mathbf{X}^s(t) \equiv (X_1^s(t), \dots, X_{N_s}^s(t))$.
 - **Criterion:** A species is *fast* if its population *gets changed* by at least one fast reaction; otherwise, the species is *slow*.

Some subtle points:

- ✓ $\nu_{ij} \rightarrow \nu_{ij}^{\sigma\rho}$, where $\sigma = f, s$ (species) and $\rho = f, s$ (reaction).
- ✓ A slow species cannot get changed by a fast reaction ($\nu_{ij}^{sf} \equiv 0$), but a fast species *might* get changed by a slow reaction.
- ✓ $a_j^s(\mathbf{x}) = a_j^s(\mathbf{x}^f, \mathbf{x}^s)$, $a_j^f(\mathbf{x}) = a_j^f(\mathbf{x}^f, \mathbf{x}^s)$.
- ✓ The population of a fast species *need not be large*.
- ✓ There might be *no slow species* (in which case $\mathbf{x}^f = \mathbf{x}$).

3rd Define the “virtual fast process” $\hat{\mathbf{X}}^f(t)$.

- $\hat{\mathbf{X}}^f(t) \equiv$ the fast species populations driven by *only* the fast reactions.
- Thus, $\hat{\mathbf{X}}^f(t)$ is $\mathbf{X}^f(t)$ with all slow reactions “turned off”.
- $\mathbf{X}^f(t)$ is a *physically real* process. But it’s *non-Markovian*.
 - Determining $\mathbf{X}^f(t)$ is no easier than determining $\mathbf{X}(t)$.
- $\hat{\mathbf{X}}^f(t)$ is a *physically fictitious* process. But it’s Markovian.
 - It satisfies an ordinary master equation – for the fast species driven by only the fast reactions (so all the slow species populations will be constant).
- The CME for $\hat{\mathbf{X}}^f(t)$ will be *simpler* than the CME for $\mathbf{X}(t)$.
It determines

$$\hat{P}(\mathbf{x}^f, t | \mathbf{x}_0^f, \mathbf{x}_0^s, t_0) \triangleq \Pr\{\hat{\mathbf{X}}^f(t) = \mathbf{x}^f | \mathbf{X}(t_0) = (\mathbf{x}_0^f, \mathbf{x}_0^s)\}.$$

4th Require the system to be “stiff”:

- **Stability Condition.** $\hat{\mathbf{X}}^f(t)$ must be *stable*; i.e., the limit

$$\lim_{t \rightarrow \infty} \hat{P}(\mathbf{x}^f, t | \mathbf{x}_0, t_0) \equiv \hat{P}(\mathbf{x}^f, \infty | \mathbf{x}_0)$$

must exist. $\hat{P}(\mathbf{x}^f, \infty | \mathbf{x}_0)$ will be the solution of

$$0 = \sum_{j=1}^{M_f} \left\{ a_j^f(\mathbf{x}^f - \mathbf{v}_j^f, \mathbf{x}_0^s) \hat{P}(\mathbf{x}^f - \mathbf{v}_j^f, \infty | \mathbf{x}_0) - a_j^f(\mathbf{x}^f, \mathbf{x}_0^s) \hat{P}(\mathbf{x}^f, \infty | \mathbf{x}_0) \right\}.$$

- **Timescale Separation Condition.** $\hat{\mathbf{X}}^f(t) \rightarrow \hat{\mathbf{X}}^f(\infty)$ in a time that is *small* compared to the *expected* time to the *next slow* reaction.

- If these conditions are satisfied, the fast reactions should be “less important” than the slow ones. Skipping over them should be okay.
- If these conditions are not satisfied, *and cannot be satisfied by any repartitioning*, then the fast reactions are no less important than the slow ones. Skipping over them is not a good idea.

5th The Slow-Scale Approximation

Lemma: With $\mathbf{X}(t) = (\mathbf{x}^f, \mathbf{x}^s)$, let Δ_s be a time increment that is *large* compared to the time for $\hat{\mathbf{X}}^f(t \rightarrow \infty) \rightarrow \hat{\mathbf{X}}^f(\infty)$, but *small* compared to the expected time to the next *slow* reaction. Then the probability that R_j^s will fire in $[t, t + \Delta_s)$ is *approximately* $\Delta_s \times \bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s)$, where

$$\bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s) \equiv \left\langle a_j^s(\hat{\mathbf{X}}^f(\infty), \mathbf{x}^s) \right\rangle = \sum_{\mathbf{y}^f} \hat{P}(\mathbf{y}^f, \infty | \mathbf{x}^f, \mathbf{x}^s) a_j^s(\mathbf{y}^f, \mathbf{x}^s).$$

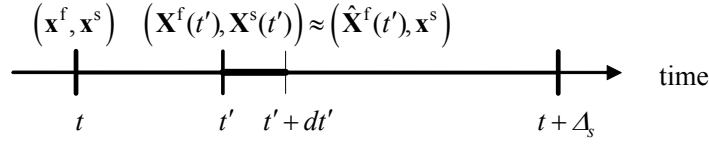
Why this lemma does the trick for us:

- For a “true” dt , $a_j^s(\mathbf{x}^f, \mathbf{x}^s) \times dt = \text{Prob}\{R_j^s \text{ in } [t, t + dt)\}$.

This is the *defining property* of the propensity function!

- Δ_s , though large on the “fast” scale, is *very small* on the “slow” scale, and the Lemma $\Rightarrow \bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s) \times \Delta_s \approx \text{Prob}\{R_j^s \text{ in } [t, t + \Delta_s)\}$.
- $\therefore \bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s)$ can be regarded as the **effective propensity function** of R_j^s on the time scale of the **slow** reactions.

Proof of the lemma:



$$\text{Prob}\{R_j^s \text{ in } [t', t' + dt']\} = a_j^s(\mathbf{X}^f(t'), \mathbf{X}^s(t')) dt' \approx a_j^s(\hat{\mathbf{X}}^f(t'), \mathbf{x}^s) dt'.$$

$$\begin{aligned} \text{Prob}\{R_j^s \text{ in } [t, t + \Delta_s]\} &\approx \int_{t'=t}^{t'+\Delta_s} a_j^s(\hat{\mathbf{X}}^f(t'), \mathbf{x}^s) dt' \\ &= \left\{ \frac{1}{\Delta_s} \int_t^{t+\Delta_s} a_j^s(\hat{\mathbf{X}}^f(t'), \mathbf{x}^s) dt' \right\} \Delta_s \\ &\approx \left\{ \text{temporal average of } a_j^s(\hat{\mathbf{X}}^f(t'), \mathbf{x}^s) \right\} \Delta_s \\ &\approx \left\{ \text{ensemble average of } a_j^s(\hat{\mathbf{X}}^f(\infty), \mathbf{x}^s) \right\} \Delta_s \\ &= \left\{ \sum_{\mathbf{y}^f} \hat{P}(\mathbf{y}^f, \infty | \mathbf{x}^f, \mathbf{x}^s) a_j^s(\mathbf{y}^f, \mathbf{x}^s) \right\} \Delta_s. \quad \text{QED} \end{aligned}$$

- The Slow-Scale SSA -

- In $(\mathbf{x}^f, \mathbf{x}^s)$ at t , evaluate $\bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s) = \langle a_j^s(\hat{\mathbf{X}}^f(\infty), \mathbf{x}^s) \rangle$, $j = 1, \dots, M_s$.
This is the hard part. Approximations may be necessary to get the first and second moments of $\hat{\mathbf{X}}^f(\infty)$; e.g., $\langle \hat{\mathbf{X}}^f(\infty) \rangle \approx \hat{\mathbf{X}}^{\text{RRE}}(\infty)$.
- With $\bar{a}_0^s(\mathbf{x}^f, \mathbf{x}^s) \equiv \sum_{j=1}^{M_s} \bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s)$,
 - take the time τ to the *next slow* reaction to be the exponential random variable with mean $1/\bar{a}_0^s(\mathbf{x}^f, \mathbf{x}^s)$;
 - take the index j of that reaction to be the integer random variable with probability $\bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s)/\bar{a}_0^s(\mathbf{x}^f, \mathbf{x}^s)$.
- Implement the next slow reaction by *first* updating

$$t \leftarrow t + \tau \quad \text{and} \quad \mathbf{x} \leftarrow \mathbf{x} + \boldsymbol{\nu}_j^s,$$
 and *then* “relaxing” the *fast* variables,

$$\mathbf{x}^f \leftarrow \text{sample of } \hat{\mathbf{X}}^f(\infty).$$

- A Variation -

Recall in our proof of the Lemma that

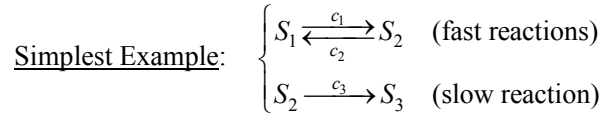
$$\bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s) = \left\langle a_j^s(\hat{\mathbf{X}}^f(\infty), \mathbf{x}^s) \right\rangle \approx \frac{1}{\Delta_s} \int_t^{t+\Delta_s} a_j^s(\hat{\mathbf{X}}^f(t'), \mathbf{x}^s) dt'$$

➤ If the *ensemble average* cannot be conveniently and accurately calculated, then compute the *temporal average* instead:

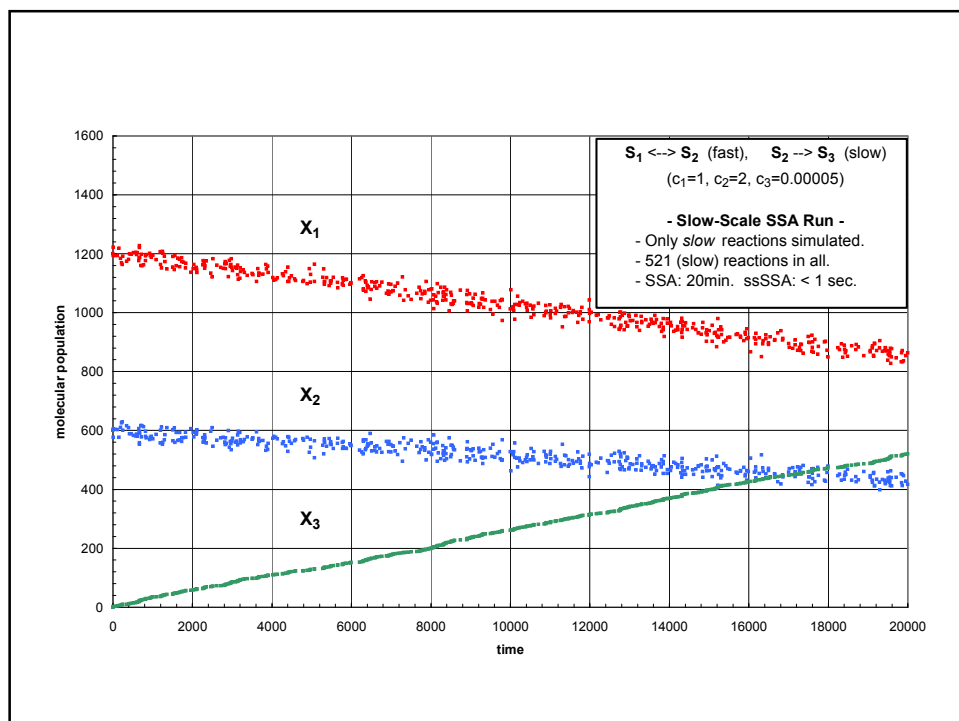
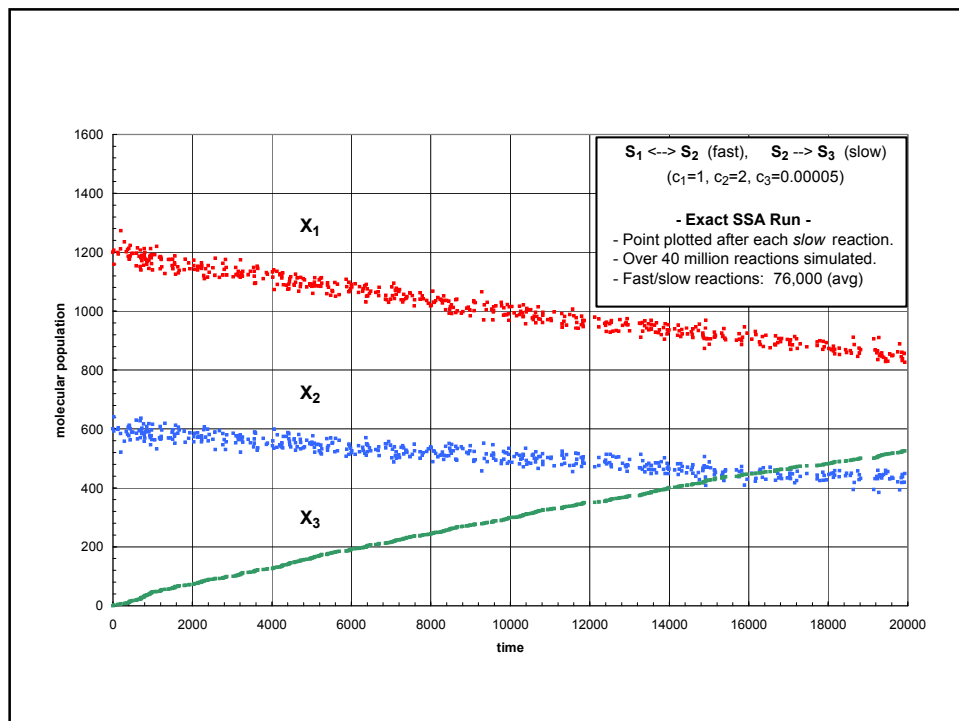
- Use the SSA to *simulate* $\hat{\mathbf{X}}^f(t')$ from time t to time $t + T_f$, where T_f is “large” compared to the relaxation time of $\hat{\mathbf{X}}^f(t')$.
- Numerically evaluate $\int_t^{t+T_f} a_j^s(\hat{\mathbf{X}}^f(t'), \mathbf{x}^s) dt' / T_f \approx \bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s)$.
- Take the random sample of $\hat{\mathbf{X}}^f(\infty)$ to be $\hat{\mathbf{X}}^f(t + T_f)$.

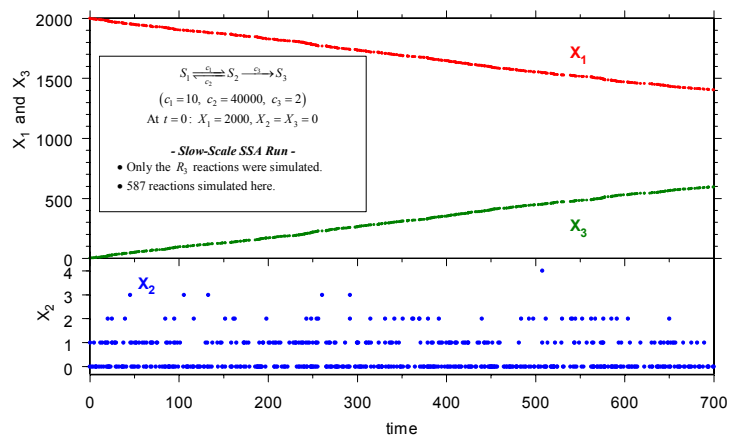
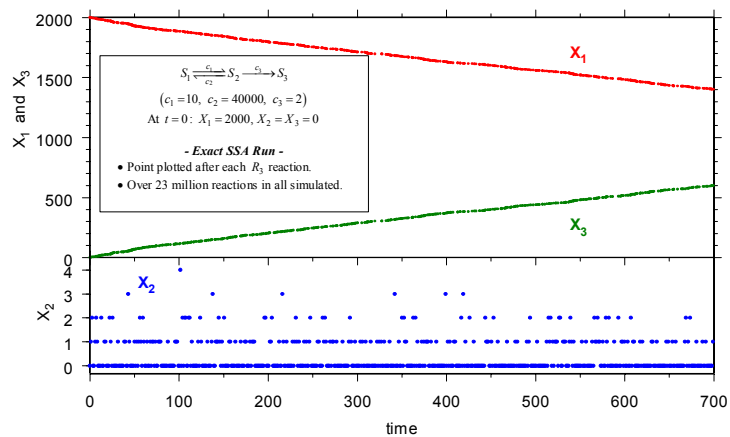
This is the “nested SSA” of ...

- E W, Liu D, Vanden-Eijnden E. 2005. Nested stochastic simulation algorithm for chemical kinetic systems with disparate rates. *J. Chem. Phys.* **123**:194107.

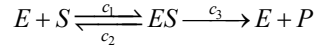


- $\mathbf{X}^f(t) = (X_1(t), X_2(t))$, $\mathbf{X}^s(t) = X_3(t)$.
- The *virtual fast process* is $\hat{\mathbf{X}}^f(t) = (\hat{X}_1(t), \hat{X}_2(t))$ for $S_1 \xrightleftharpoons[c_2]{c_1} S_2$.
 - $\hat{X}_1(t) + \hat{X}_2(t) = x_{12}$. $\hat{\mathbf{X}}^f(\infty)$ can be calculated exactly:
 - $\hat{X}_2(\infty) = \text{Bin}\left(\frac{c_1}{c_1 + c_2}, x_{12}\right)$, $\hat{X}_1(\infty) = x_{12} - \hat{X}_2(\infty)$.
- It follows that: $\bar{a}_3(\mathbf{x}) \equiv c_3 \left\langle \hat{X}_2(\infty) \right\rangle = \frac{c_3 c_1 x_{12}}{c_1 + c_2}$.
- The *relaxation time* for $\hat{X}^f(t) \rightarrow \hat{X}^f(\infty)$ is $1/(c_1 + c_2)$; thus, the condition for using the ssSSA is: $1/(c_1 + c_2) \ll (c_1 + c_2)/c_3 c_1 x_{12}$. This condition will be satisfied if c_2 is sufficiently $\gg c_3$.





The Enzyme-Substrate Reaction



- Fast reactions: R_1 and R_2 . Slow reaction: R_3 .
- Fast species: $\mathbf{X}^f = (X_E, X_S, X_{ES})$. Slow species: $\mathbf{X}^s = X_P$.
- Virtual fast process: $\hat{\mathbf{X}}^f = (\hat{X}_E, \hat{X}_S, \hat{X}_{ES})$ for $E + S \xrightleftharpoons[c_2]{c_1} ES$.

Two conservation eqns: $\hat{X}_E(t) + \hat{X}_{ES}(t) = x_{E*}$, $\hat{X}_S(t) + \hat{X}_{ES}(t) = x_{S*}$.
 \Rightarrow one independent variable, say $\hat{X}_{ES}(t)$. For it we have:

$$\begin{aligned} a_1(x_{ES}) &= c_1(x_{E*} - x_{ES})(x_{S*} - x_{ES}), & v_1 &= +1; \\ a_2(x_{ES}) &= c_2 x_{ES}, & v_2 &= -1. \end{aligned}$$

The resulting $t \rightarrow \infty$ master equation for $\hat{X}_{ES}(t)$ implies the “detailed balance condition”.

$$\begin{aligned} c_2(x_{ES} + 1)\hat{P}(x_{ES} + 1, \infty | x_{E*}, x_{S*}) \\ = c_1(x_{E*} - x_{ES})(x_{S*} - x_{ES})\hat{P}(x_{ES}, \infty | x_{E*}, x_{S*}), \quad \forall x_{ES} \\ \dots \text{ which gives us a recursion relation for } \hat{P}(x_{ES}, \infty | x_{E*}, x_{S*}): \end{aligned}$$

$$\begin{aligned} \bullet \quad \hat{P}(x_{ES} + 1, \infty | x_{E*}, x_{S*}) &= \frac{c_1(x_{E*} - x_{ES})(x_{S*} - x_{ES})}{c_2(x_{ES} + 1)} \hat{P}(x_{ES}, \infty | x_{E*}, x_{S*}) \\ &\quad (x_{ES} = 0, 1, \dots, x_{ES}^{\max} - 1) \end{aligned}$$

where $x_{ES}^{\max} \equiv \min(x_{E*}, x_{S*})$.

- In principle, we can compute $\hat{P}(x_{ES}, \infty | x_{E*}, x_{S*})$ from this recursion relation *exactly*. Then we can *exactly*:

$$\text{- compute } \bar{a}_3(\mathbf{x}) = c_3 \left\langle \hat{X}_{ES}(\infty) \right\rangle = c_3 \sum_{x_{ES}=0}^{x_{ES}^{\max}} x_{ES} \hat{P}(x_{ES}, \infty | x_{E*}, x_{S*});$$

$$\text{- generate } \hat{X}_{ES}(\infty) \text{ by the rule } \sum_{x_{ES}=0}^{\hat{X}_{ES}(\infty)} \hat{P}(x_{ES}, \infty | x_{E*}, x_{S*}) \geq r,$$

$$\text{and then get } \hat{X}_E(\infty) = x_{E*} - \hat{X}_{ES}(\infty) \text{ and } \hat{X}_S(\infty) = x_{S*} - \hat{X}_{ES}(\infty).$$

- For “small” x_{ES}^{\max} , all this is feasible. So we do it.
- For “large” x_{ES}^{\max} , we have to make some *approximations*.

- **Estimating $\bar{a}_3(\mathbf{x}) = c_3 \langle \hat{X}_{\text{ES}}(\infty) \rangle$ for “large” $x_{\text{ES}}^{\text{max}}$** -

- Tests show that in this case, $\langle \hat{X}_{\text{ES}}(\infty) \rangle \approx \bar{x}_{\text{ES}}$, where:

$$0 = c_1(x_{\text{E}^*} - \bar{x}_{\text{ES}})(x_{\text{S}^*} - \bar{x}_{\text{ES}}) - c_2\bar{x}_{\text{ES}}.$$

Solving this gives

$$\bar{x}_{\text{ES}} = \frac{1}{2} \left\{ \left(x_{\text{E}^*} + x_{\text{S}^*} + \frac{c_2}{c_1} \right) - \sqrt{\left(x_{\text{E}^*} + x_{\text{S}^*} + \frac{c_2}{c_1} \right)^2 - 4x_{\text{E}^*}x_{\text{S}^*}} \right\}.$$

- So for large $x_{\text{ES}}^{\text{max}}$, we approximate $\bar{a}_3(\mathbf{x}) \approx c_3\bar{x}_{\text{ES}}$.

➤ **Connection to Michaelis-Menten:** Can prove that

$$c_3\bar{x}_{\text{ES}} = \frac{c_3x_{\text{E}^*}\bar{x}_{\text{S}}}{(c_2/c_1) + \bar{x}_{\text{S}}} (\equiv v_{\text{MM}})$$

when $\bar{x}_{\text{S}} = x_{\text{S}^*} - \bar{x}_{\text{ES}}$ and we invoke the **rapid equilibrium assumption**

$$c_1(\bar{x}_{\text{E}^*} - \bar{x}_{\text{ES}})\bar{x}_{\text{S}} = c_2\bar{x}_{\text{ES}}.$$

We have thus **derived** the Michaelis-Menten formula as a *large population approximation to $\bar{a}_3(\mathbf{x})$* , but *without* having to make any *assumption* of “rapid equilibrium” or “quasi-steady state”.

- **Generating a random sample of $\hat{X}_{\text{ES}}(\infty)$ for “large” $x_{\text{ES}}^{\text{max}}$** -

- For this, tests show that we can approximate $\hat{X}_{\text{ES}}(\infty)$ as a *normal* random variable whose *mean* is \bar{x}_{ES} ...

... and whose *variance* is the variance $\bar{\sigma}_{\text{G}}^2$ of the Gaussian curve that best fits the *peak* of the function $\hat{P}(x_{\text{ES}}, \infty | x_{\text{E}^*}, x_{\text{S}^*})$ defined by the recursion relation. Can show that

$$\bar{\sigma}_{\text{G}}^2 = \frac{c_2\tilde{x}_{\text{ES}}}{-2c_1\tilde{x}_{\text{ES}} + c_1(x_{\text{E}^*} + x_{\text{S}^*} + 2) + c_2},$$

where

$$\tilde{x}_{\text{ES}} = \frac{1}{2} \left\{ \left(x_{\text{E}^*} + x_{\text{S}^*} + \frac{c_2}{c_1} + 2 \right) - \sqrt{\left(x_{\text{E}^*} + x_{\text{S}^*} + \frac{c_2}{c_1} + 2 \right)^2 - 4(x_{\text{E}^*} + 1)(x_{\text{S}^*} + 1)} \right\}.$$

- So we get $\hat{X}_{\text{ES}}(\infty)$ by Monte Carlo sampling $\mathcal{N}(\bar{x}_{\text{ES}}, \bar{\sigma}_{\text{G}}^2)$. Then we take $\hat{X}_{\text{E}}(\infty) = x_{\text{E}^*} - \hat{X}_{\text{ES}}(\infty)$ and $\hat{X}_{\text{S}}(\infty) = x_{\text{S}^*} - \hat{X}_{\text{ES}}(\infty)$.

➤ **Note:** Any errors we make in generating the fast species populations will **not** affect the accuracy of the slow reaction simulations!

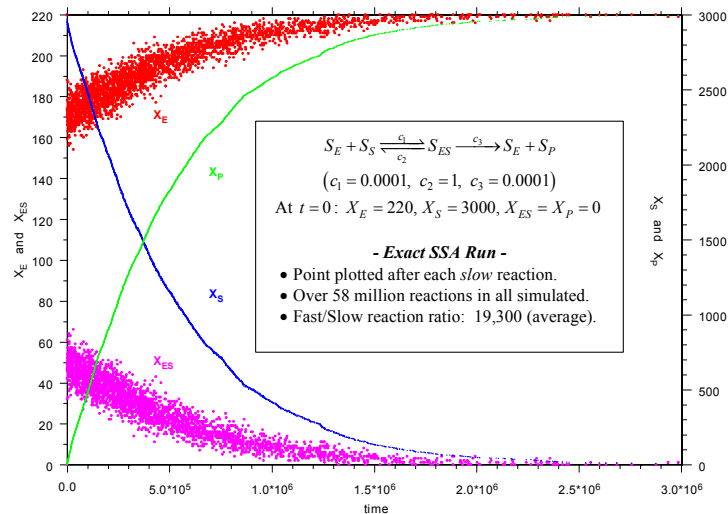
- The **stochastic stiffness condition** for this reaction set is:

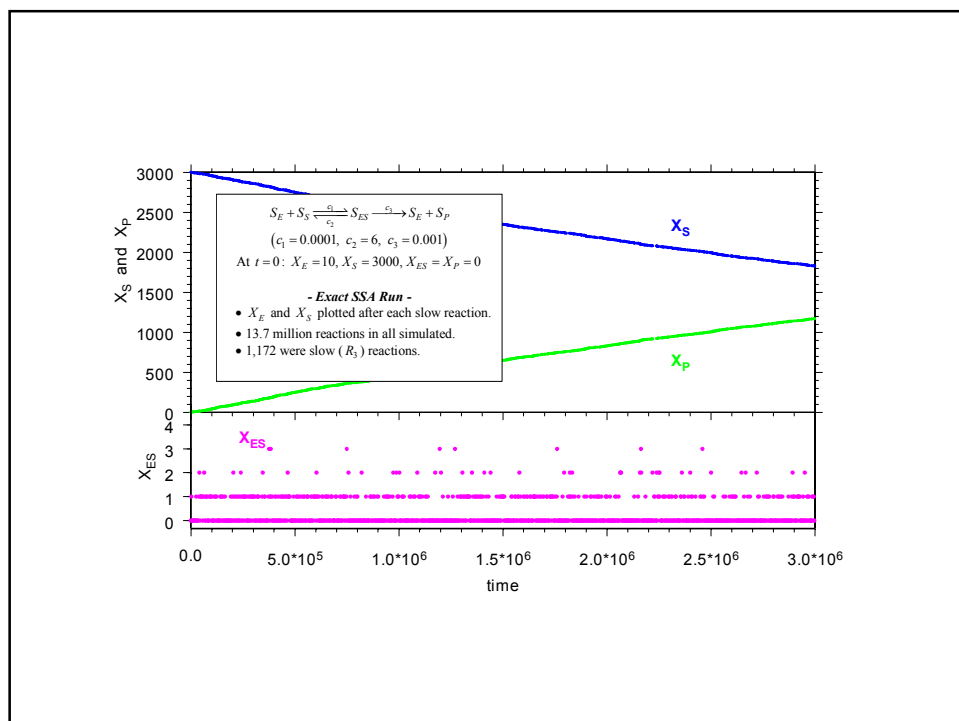
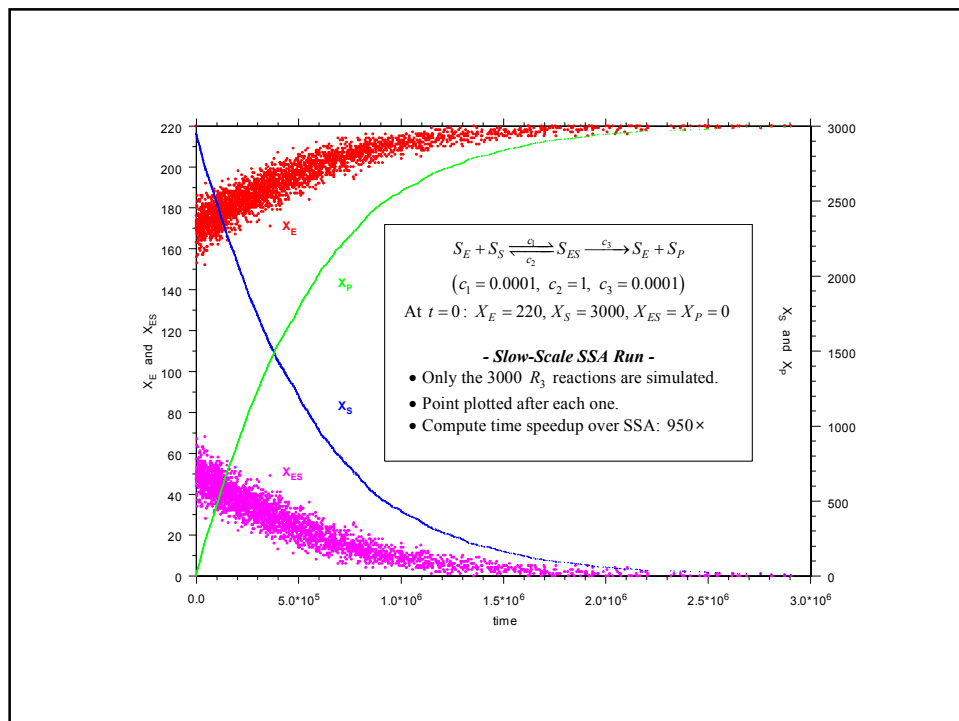
$\{\text{relaxation time of } \hat{X}_{ES}(t)\} \ll \{\text{expected time to the next } R_3 \text{ reaction}\}.$

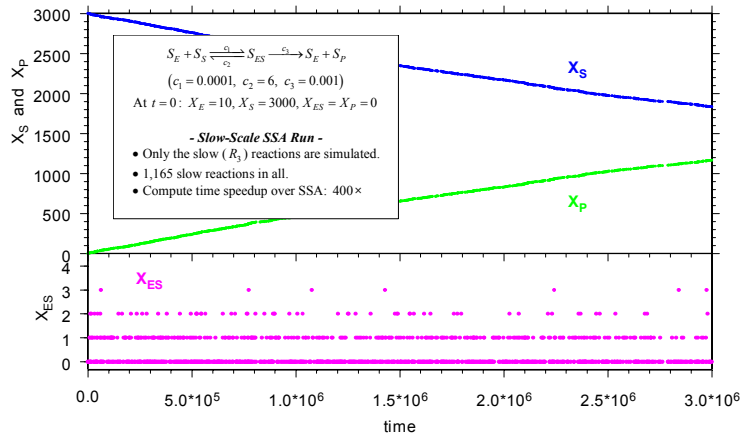
- We can expect the ssSSA to give reliable results **if and only if** this condition is satisfied.
- This condition turns out to be:

$$\frac{1}{-2c_1\tilde{x}_{ES} + c_1(x_{E*} + x_{S*} + 2) + c_2} \ll \frac{1}{c_3\tilde{x}_{ES}}.$$

- It will be satisfied if c_2 is sufficiently larger than c_3 .
- **Physical Interpretation:** When $c_2 \gg c_3$, ES will decay to $E + S$ much more often than to $E + P$. \Rightarrow Between two *successive* R_3 reactions there will be *very many* R_1 and R_2 reactions. $\Rightarrow R_1$ and R_2 will “equilibrate” before the next R_3 reaction occurs. \Rightarrow Lemma.







Summary: The ssSSA ...

- Provides a detailed operational procedure for determining
 - the fast and slow **reactions**,
 - the fast and slow **species**,
 - a **virtual fast process** $\hat{\mathbf{X}}^f(t)$ that is **tractable**.
- Requires the **stiffness conditions** to be satisfied:
 - $\hat{\mathbf{X}}^f(\infty)$ must *exist*;
 - $\hat{\mathbf{X}}^f(t) \rightarrow \hat{\mathbf{X}}^f(\infty)$ *rapidly* on the time scale of the *slow* reactions.
- Is based on the **Slow-Scale Approximation Lemma**:

On the time scale of the **slow** reactions, $a_j^s(\mathbf{x}^f, \mathbf{x}^s)$ can be **approximately replaced** by $\bar{a}_j^s(\mathbf{x}^f, \mathbf{x}^s) \triangleq \langle a_j^s(\hat{\mathbf{X}}^f(\infty), \mathbf{x}^s) \rangle$.
- Simulates *only the slow reactions*, skipping over the numerous and uninteresting fast reactions, yet accurately replicates the trajectories of *all* the species.