

Modelling and Verification of Spatiality: The Shape Calculus

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OPTC Workshop - Bertinoro, 18 - 21 June 2014

Explicitly Spatial Models and Languages

- For modeling Distributed Systems with a known spatial distribution and/or moving components, e.g. protocols for mobile systems with positions/trajectories
- For modelling Biological Systems:
 - **Spatiality**: proximity, affinity, crowding, co-localization, perception, . . .
 - **Shape/Geometry**: the shape usually determines the function
 - **Movement**: biological entities can move with different patterns (Brownian motion, chemical gradient-guided, . . .)

Which Space Representation?

- Topological space
 - Compartments, no explicit coordinates - Interactions occur among entities in the same compartment
 - e.g. Membrane Computing (1998), Bio Ambients (2004), Brane Calculi (2004)
- Discrete Grid space
 - 2D or 3D cells, discrete coordinates - Interactions by proximity of cells, discrete movements
 - e.g. Cellular Automata (1940s), Spatial P Systems (2010)
- Continuous Geometrical space
 - 3D positions, with continuous coordinates - Interactions can be defined in different ways, using geometrical information
 - e.g. Real-space Process Algebra (Spatio-temporal ACP, 1993!), Space- π (2008), 3π (2012), Shape Calculus (2010)

*Examples mainly Bio-oriented languages

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Challenging Features

The area is relatively new, there is a lot of room for working :-)

- Representation of Space and Time
- Expression of (Complex) Movements
- Space-driven Interactions vs Classical communication
- Shapes, Bonds, Re-shaping, Growth, Reversibility of Reactions
- How to deal with collisions
- Modelling and Simulation, what about Verification?

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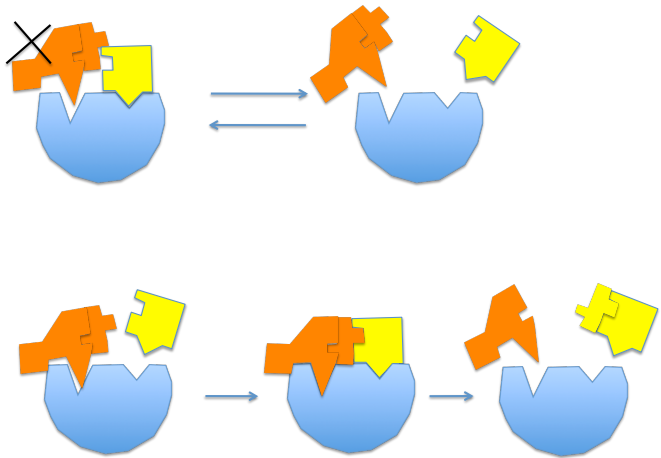
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What we have done so far...

- We defined a process calculus with a very rich set of features: the Shape Calculus
- Shape Calculus models are stub code for entities in a simulator
- We retain the possibility of formal **verification**, e.g. by abstract interpretation
- We can combine **simulation** (with a less effort of coding) with **verification**

Scenario we want to model

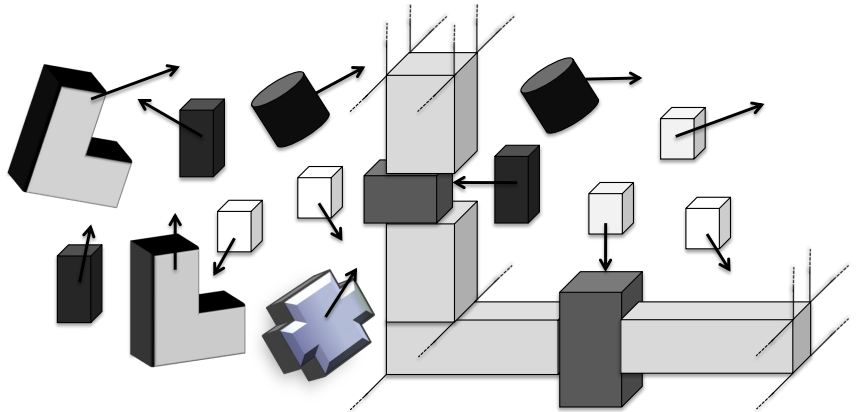


- Spatial 3D, Deterministically Timed, Individual-based
- Shape-based \rightarrow modelled entities have a 3D **geometrical** shape, a position and a velocity with which they move
- Entities can interact \rightarrow bindings (collision-driven) and splittings (driven by internal behaviour)

[†]E. Bartocci, F. Corradini, M. R. Di Berardini, E. Merelli, and L. Tesei, *Shape Calculus. A Spatial Mobile Calculus for 3D Shapes*, Scientific Annals of Computer Science **20** (2010), pp. 1-31.

[†]E. Bartocci, D. R. Cacciagrano, M. R. Di Berardini, E. Merelli, and L. Tesei, *Shape Calculus: Timed Operational Semantics and Well-Formedness*, Scientific Annals of Computer Science **20** (2010), pp. 32-52.

Shape Calculus: Scenario



What is a 3D spatial process, geometrically?

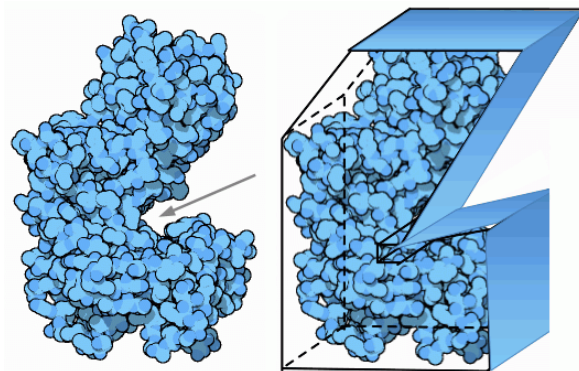
In most of the languages with explicit space, a process is a point $\vec{p} \in \mathbb{R}^3$

- Real space, but no dimension
- Interactions can depend only on distance (e.g. in Space- π or in spatio-temporal ACP)

In Shape Calculus, a process is a set of points $V \subseteq \mathbb{R}^3$

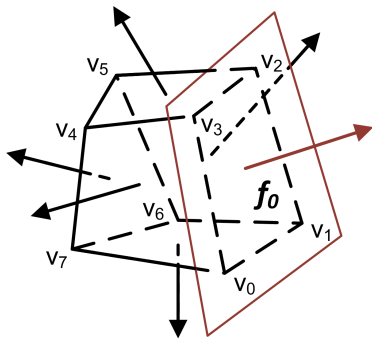
- A “physical body” is associated to the process, its **shape**
- Interactions can depend on distance, but also on the shape and on the 3D orientation of the process

Shape Approximation



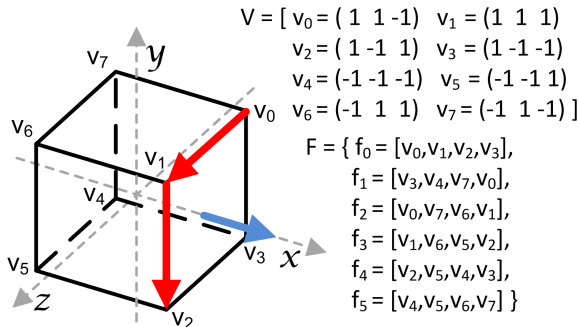
Basic Shapes, Shape Calculus Approach

- We take only *convex polyhedra* as basic shapes
- A convex polyhedron is determined by the planes of its faces and their outward normals \rightarrow it is defined by its *features*



Basic Shapes

- We are interested in the features, which can be communication channels, more than in the internal points
- Basic shapes are represented as *Indexed face set*, i.e. as point tuples (faces)



Shape Syntax

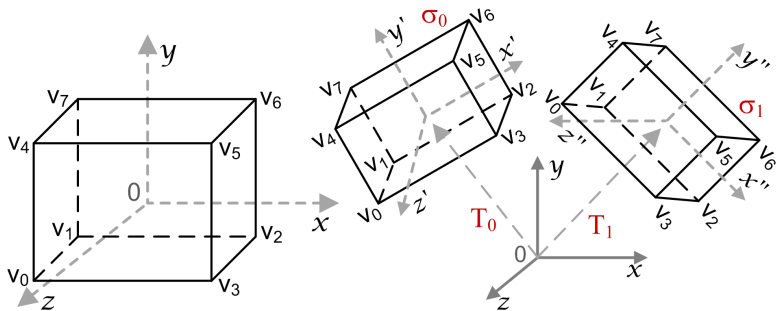
$$S ::= \sigma \mid S \langle \varphi_1, \varphi_2 \rangle S$$

$\sigma \in \text{Basic}$, φ_1, φ_2 are two features, one for each of the shapes, W is an IFS, ω is the angular velocity (quaternion) and s is a **name** identifying univocally the shape

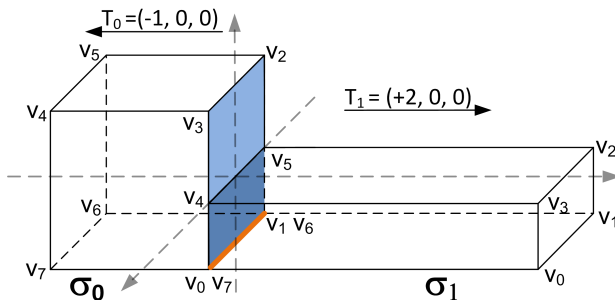
$$\sigma = \langle W, m, \mathbf{v}, \omega, s, T \rangle$$

- W is the *template* shape, with coordinates expressed in the *local frame*
- W is moved to the *reference frame* by the **affine transformation** T associated to the shape (similar to 3π)

IFSs and shapes

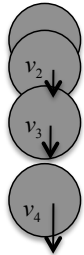


Compound Shapes

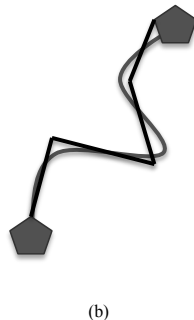


- Two faces touch: $s_0.f_0 = [s_0.v_0, s_0.v_1, s_0.v_2, s_0.v_3]$,
 $s_1.f_1 = [s_1.v_4, s_1.v_5, s_1.v_6, s_1.v_7]$
- The compound shape is $S_1 = \sigma_0 \langle \{s_0.f_0, s_1.f_1\} \rangle \sigma_1$

Movement


$$\text{Move}(S, t_i) = [0, -\frac{1}{2}g(i+1)\Delta, 0]m/s$$
$$t_i = t_{i-1} + \Delta$$
$$\Delta = 0.05s$$
$$v_1 = \text{Move}(S, t_0) = [0, -0.245, 0]m/s$$
$$v_2 = \text{Move}(S, t_1) = [0, -0.49, 0]m/s$$
$$v_3 = \text{Move}(S, t_2) = [0, -0.735, 0]m/s$$
$$v_4 = \text{Move}(S, t_3) = [0, -0.98, 0]m/s$$
$$S = \text{circle}$$
$$v = [v_x, v_y, v_z]$$

(a)



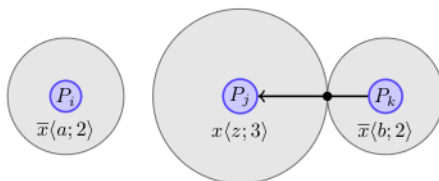
Time evolution and velocity update

- Time domain is continuous, divided into **small** discrete time steps Δ
- At each step the velocity and rotation of each entity is changed by the Move function

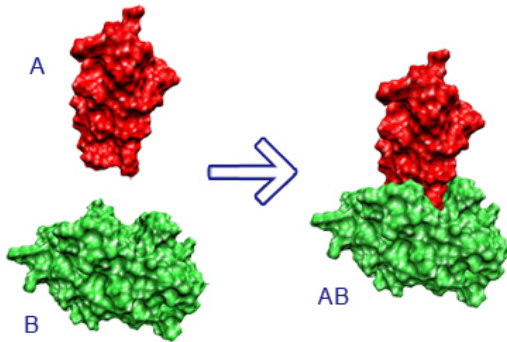
Spatial-dependent communication

In languages with no shapes, communication:

- Can depend on distance
- Can be synchronous, checking distances at regular time steps (Space- π)
- Can be asynchronous, with spherical propagation of messages (Spatio-temporal ACP)
- Classical communication can be re-obtained, allowing infinite distance

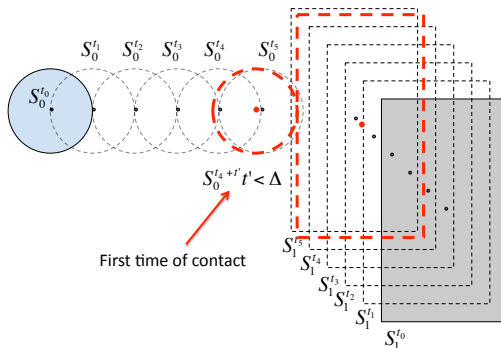


Interactions of 3D processes: binding and splitting



In Shape Calculus, Communication is synchronous and only local (collision-driven)

Collision Detection



- FTOC can be determined in each Δ interval, if any
- In case of collision the time delay step is reduced from Δ to FTOC
- Collision(s) are detected and may trigger **interactions**

Behaviour Syntax

$$B ::= \text{nil} \mid \langle \alpha, \{\varphi_1, \dots, \varphi_n\} \rangle . B \mid \epsilon(t).B \mid B + B \mid K \mid$$
$$\omega(a, \{\varphi_1, \dots, \varphi_n\}).B \mid \rho(L).B$$

where $\varphi_1, \dots, \varphi_n$ are feature representations and
 $L = \{ \langle a, \{\varphi_1, \dots, \varphi_n\} \rangle \mid \langle a, \{\varphi_1, \dots, \varphi_n\} \rangle \text{ is a channel} \}$.

- **Spatial channels** can be open on several features of the shape, even disjoint
- Communication is possible only between compatible spatial channels that are **colliding**

SOS rules for Behaviours

$$\text{PREF}_a \frac{\mu \in \mathcal{C} \cup \omega(\mathcal{C})}{\mu.B \xrightarrow{\mu} B}$$

$$\text{DEL}_a \frac{B \xrightarrow{\mu} B'}{\epsilon(0).B \xrightarrow{\mu} B'}$$

$$\text{SUM}_a \frac{B_1 \xrightarrow{\mu} B'_1}{B_1 + B_2 \xrightarrow{\mu} B'_1}$$

$$\text{STR}_1 \frac{L = \{\langle \alpha, X \rangle\}}{\rho(L).B \xrightarrow{\rho(\alpha, X)} B}$$

$$\text{STR}_2 \frac{L = \{\langle \alpha, X \rangle\} \cup L' \quad L' \neq \emptyset}{\rho(L).B \xrightarrow{\rho(\alpha, X)} \rho(L').B}$$

$$\text{STR}_3 \frac{B \xrightarrow{\rho(\alpha, X)} B'}{\rho(L).B \xrightarrow{\rho(\alpha, X)} \rho(L).B'}$$

$$\text{NIL}_t \frac{}{\text{nil} \rightsquigarrow \text{nil}}$$

$$\text{PREF}_t \frac{\mu \in \mathcal{C} \cup \omega(\mathcal{C})}{\mu.B \rightsquigarrow \mu.B}$$

$$\text{STR}_t \frac{}{\rho(L).B \rightsquigarrow \rho(L).B}$$

$$\text{SUM}_t \frac{B_1 \rightsquigarrow B'_1 \quad B_2 \rightsquigarrow B'_2}{B_1 + B_2 \rightsquigarrow B'_1 + B'_2}$$

$$\text{DEL}_t \frac{t' \geq t}{\epsilon(t').B \rightsquigarrow \epsilon(t' - t).B}$$

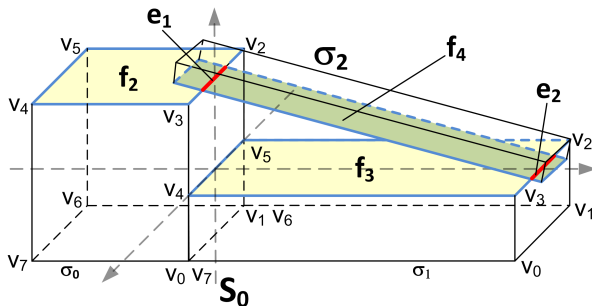
$$\text{DEL}_c \frac{B \rightsquigarrow B'}{\epsilon(t').B \rightsquigarrow^{t+t'} B'}$$

3D Process Syntax

$$P ::= S[B] \mid P \langle a, \{\varphi_1, \varphi_2\} \rangle P$$

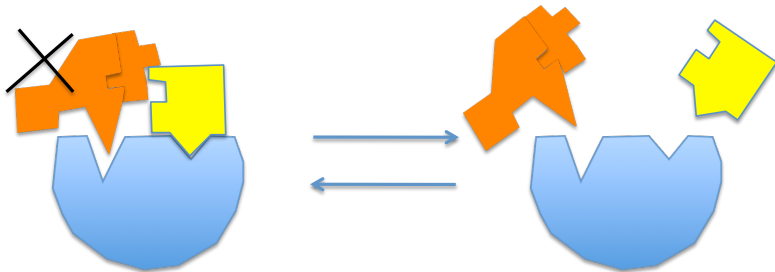
where a is the name of the channel and φ_1, φ_2 are features belonging to the shapes associated with the compound 3D processes.

- **Composition by glueing**, after a collision on compatible active channels on surfaces
- Composed process behaviour is the **interleaving** of the components' behaviours



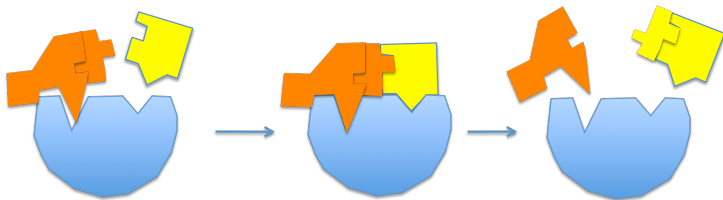
- Non-deterministically, one of the two possible bonds is chosen (s_2 is the name of σ_2):
 - $S_0[B_0] \langle a, \{s_0.e_1, s_2.f_4\} \rangle \sigma_2[B_1]$
 - $S_0[B_0] \langle a, \{s_1.e_2, s_2.f_4\} \rangle \sigma_2[B_1]$

Weak Split - $\omega(C)$



A bond on a channel C can be released non-deterministically

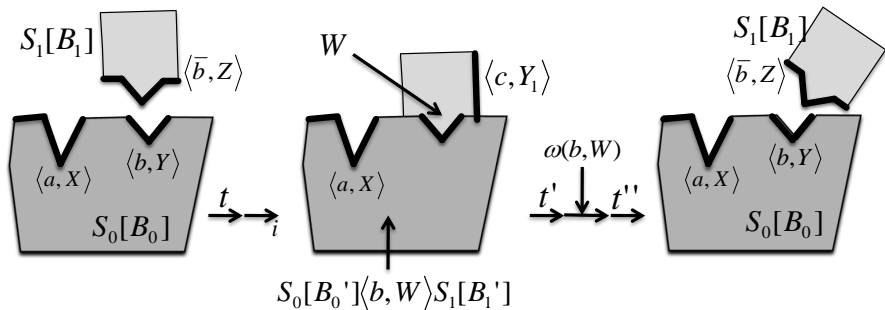
Strong Splitting $\rho(L)$



A subset L of bonds must split at the same time

A “re-shaping” is obtained by rearrangement of bonds

Binding and Splitting



SOS rules for 3D processes

$$\begin{array}{c}
 \text{BASIC}_t \frac{B \xrightarrow{t} B'}{S[B] \xrightarrow{t} (S + t)[B']} \quad \text{COMP}_t \frac{P \xrightarrow{t} P' \quad Q \xrightarrow{t} Q' \quad X' = X + (t \cdot \mathbf{v}(P))}{P \langle a, X \rangle Q \xrightarrow{t} P' \langle a, X' \rangle Q'} \\
 \\
 \text{BASIC}_c \frac{B \xrightarrow{\langle \alpha, X \rangle} B' \quad Y = \text{global}(X, \mathcal{R}(S))}{S[B] \xrightarrow{\langle \alpha, Y \rangle} S[B']} \\
 \\
 \text{BASIC}_s \frac{B \xrightarrow{\rho(\alpha, X)} B' \quad Y = \text{global}(X, \mathcal{R}(S))}{S[B] \xrightarrow{\rho(\alpha, Y)} S[B']} \\
 \\
 \text{COMP}_s \frac{P \xrightarrow{\rho(\alpha, Y)} P'}{P \langle a, X \rangle Q \xrightarrow{\rho(\alpha, Y)} P' \langle a, X \rangle Q} \\
 \\
 \text{COMP}_c \frac{P \xrightarrow{\langle \alpha, Y \rangle} P' \quad Y \subseteq \mathcal{B}(P \langle a, X \rangle Q)}{P \langle a, X \rangle Q \xrightarrow{\langle \alpha, Y \rangle} P' \langle a, X \rangle Q}
 \end{array}$$

Network of 3D processes

The set of networks of 3D processes is generated by:

$$N ::= \text{Nil} \mid P \mid N \parallel N$$

Functional and Temporal semantics of the networks

$$\begin{array}{c} \text{EMPTY}_t \frac{}{\text{Nil} \xrightarrow{t} \text{Nil}} \quad \text{PAR}_t \frac{N \xrightarrow{t} N' \quad M \xrightarrow{t} M'}{N \parallel M \xrightarrow{t} N' \parallel M'} \quad \text{PAR}_{a1} \frac{N \xrightarrow{\nu} N'}{N \parallel N \xrightarrow{\nu} N' \parallel M} \end{array}$$

- $P \xrightarrow{\nu} Q$ for $\nu \in \{\omega, \rho\}$
- $P \xrightarrow{\kappa} Q$ for every collision response
- $P \xrightarrow{t} Q$ iff $P \xrightarrow{t} Q$ and either $P \not\xrightarrow{\rho}$ or $P \not\xrightarrow{\omega}$

Simulation and Verification

- The evolution of a given network can be **simulated**
- Algorithms and data structures can be imported from computational geometry
- Sources of randomness: movement function (e.g. Brownian motion)
- Sources of non-determinism: alternative choices among spatial channels, weak splits, multiple collision response
- Verification?

Verification by Abstract Interpretation

- Approach: 3D Networks are abstracted and abstract fixpoint semantics is executed
- The abstract domain and the abstract operators should guarantee termination
- First (huge) abstraction done on movement, shape and time
- Untimed and Unspatial safety properties can be checked!

†Federico Buti, Massimo Callisto De Donato, Flavio Corradini, Maria Rita Di Berardini, Emanuela Merelli, Luca Tesei: *Towards Abstraction-Based Verification of Shape Calculus*. Electr. Notes Theor. Comput. Sci. **284**: 23-34 (2012)

Concrete Set and abstraction

- $(\wp(\mathbb{N}), \subseteq, \cup, \cap, \{\}, \mathbb{N})$
- 3D processes “instances” *collapse* to a single abstract process
- A network is abstracted to a set of abstract processes
- $\epsilon(t)$ delays are collapsed to $\epsilon(\cdot)$, bindings possible at any time

Abstraction Function

| | | |
|--------------------------------------|-----|--|
| $\alpha_S(\sigma)$ | $=$ | $\langle V, m, \odot, \odot \rangle$ |
| $\alpha_S(S \langle X \rangle S)$ | $=$ | $\alpha_S(S) \langle \odot \rangle \alpha_S(S)$ |
| $\alpha_B(\text{nil})$ | $=$ | nil^\sharp |
| $\alpha_B(\langle a, X \rangle . B)$ | $=$ | $\langle a, X \rangle . \alpha_B(B)$ |
| $\alpha_B(B_1 + B_2)$ | $=$ | $\alpha_B(B_1) + \alpha_B(B_2)$ |
| $\alpha_B(\omega(a, X) . B)$ | $=$ | $\omega(a, X) . \alpha_B(B)$ |
| $\alpha_B(\rho(a, X) . B)$ | $=$ | $\rho(a, X) . \alpha_B(B)$ |
| $\alpha_B(\epsilon(t) . B)$ | $=$ | $\epsilon(\cdot) . \alpha_B(B)$ |
| $\alpha_B(\rho(L) . B)$ | $=$ | $\rho(L) . \alpha_B(B)$ |
| $\alpha_P(S[B])$ | $=$ | $\alpha_S(S)[\alpha_B(B)]$ |
| $\alpha_P(P \langle a, X \rangle Q)$ | $=$ | $\alpha_P(P) \langle a, \odot \rangle \alpha_P(Q)$ |
| $\alpha_P(\ _{i \in I} P_i)$ | $=$ | $\bigcup_{i \in I} \{\alpha(P_i)\}$ |

Abstract Set and concretization

- $(\wp(3DP^\#), \subseteq, \cup, \cap, \{\}, 3DP^\#)$
- Every abstract set is concretized considering all the possible values of velocity, position and cardinality
- Velocity of shapes is limited to a *maximal velocity* v_{max} .
- From a compounded abstract shape a set of concrete compounded shapes is generated
- Each set differs in its cardinality and in the number of concrete instances of each abstract 3D process
- The cardinality can be zero

Concretization Function

| | |
|---|---|
| $\gamma_S(\sigma^\#)$ | $= \{\sigma = \langle V, m, p, v \rangle \mid p \in \mathbb{P}, 0 \leq \ v\ \leq v_{max}\}$ |
| $\gamma_S(S^\# \langle \odot \rangle T^\#)$ | $= \{S \langle Y \rangle T \mid S \in \gamma_S(S^\#), T \in \gamma_S(T^\#), \\ Y \subseteq \mathcal{B}(S) \cap \mathcal{B}(T), Y \neq \emptyset\}$ |
| $\gamma_B(\text{nil}^\#)$ | $= \{\text{nil}\}$ |
| $\gamma_B(\langle a, X \rangle . B^\#)$ | $= \{\langle a, X \rangle . B \mid B \in \gamma_B(B^\#)\}$ |
| $\gamma_B(B_1^\# + B_2^\#)$ | $= \{B_1 + B_2 \mid B_1 \in \gamma_B(B_1^\#), B_2 \in \gamma_B(B_2^\#)\}$ |
| $\gamma_B(\omega(a, X) . B^\#)$ | $= \{\omega(a, X) . B \mid B \in \gamma_B(B^\#)\}$ |
| $\gamma_B(\rho(a, X) . B^\#)$ | $= \{\rho(a, X) . B \mid B \in \gamma_B(B^\#)\}$ |
| $\gamma_B(\epsilon(\cdot) . B^\#)$ | $= \{\epsilon(t) . B \mid t \in \mathbb{T}, B \in \gamma_B(B^\#)\}$ |
| $\gamma_B(\rho(L) . B^\#)$ | $= \{\rho(L) . B \mid B \in \gamma_B(B^\#)\}$ |
| $\gamma_P(S^\# [B^\#])$ | $= \{S[B] \mid S \in \gamma_S(S^\#), B \in \gamma_B(B^\#)\}$ |
| $\gamma_P(P^\# \langle a^\#, \odot \rangle Q^\#)$ | $= \{P \langle a, X \rangle Q \mid P \in \gamma_P(P^\#), Q \in \gamma_P(Q^\#), \\ X \subseteq \mathcal{B}(\text{shape}(P)) \cap \mathcal{B}(\text{shape}(Q)), X \neq \emptyset\}$ |
| $\gamma_P(\{P_1^\#, \dots, P_n^\#\})$ | $= \wp\{N \in \mathbb{N} \mid N = (\ _{i=1}^n (\ _{P \in U} P)), U \in \wp(\gamma_P(P_i^\#)), U \text{ finite}\}$ |

Abstract semantics

$$\text{DEL}_a^\# \frac{B^\# \xrightarrow{\mu}_\# B'^\#}{\epsilon(\cdot).B^\# \xrightarrow{\mu}_\# B'^\#}$$

$$\text{DEL}_t^\# \frac{}{\epsilon(\cdot).B^\# \rightsquigarrow_\# \epsilon(\cdot).B^\#}$$

$$\text{DEL}_c^\# \frac{B^\# \rightsquigarrow_\#^t B'^\#}{\epsilon(\cdot).B^\# \rightsquigarrow_\#^{t'} B'^\#}$$

$$\text{COMP}_t^\# \frac{P^\# \rightsquigarrow_\#^t P'^\# \quad Q^\# \rightsquigarrow_\#^t Q'^\#}{P^\# \langle a, \odot \rangle Q^\# \rightsquigarrow_\#^t P'^\# \langle a, \odot \rangle Q'^\#}$$

$$\text{COMP}_c^\# \frac{P^\# \xrightarrow{\langle a, \odot \rangle}_\# P'^\#}{P \langle a, \odot \rangle Q^\# \xrightarrow{\langle a, \odot \rangle}_\# P' \langle a, \odot \rangle Q^\#}$$

$$\text{BASIC}_c^\# \frac{B^\# \xrightarrow{\langle a, X \rangle}_\# B'^\#}{S^\#[B^\#] \xrightarrow{\langle a, \odot \rangle}_\# S^\#[B'^\#]}$$

$$\text{BASIC}_s^\# \frac{B^\# \xrightarrow{\langle a, X \rangle}_\# B'^\#}{S^\#[B^\#] \xrightarrow{\rho(\alpha, \odot)}_\# S^\#[B'^\#]}$$

$$\text{STRSYNC}^\# \frac{P^\# \xrightarrow{\rho(\alpha, \odot)}_\# P'^\# \quad Q^\# \xrightarrow{\rho(\bar{\alpha}, \odot)}_\# Q'^\# \quad \alpha \in \{a, \bar{a}\}}{P^\# \langle a, \odot \rangle Q^\# \xrightarrow{\rho(a, \odot)}_\# P' \langle a, \odot \rangle Q'}$$

Verification of Safety Properties

- Given a network N we calculate the least fixpoint of the abstract semantics:
 - $F \uparrow^0 (N) = \alpha(\{N\})$
 - $F \uparrow^n (N) = \bigcup_{M^\# \in \{N^\# \mid F \uparrow^{n-1}(N) \xrightarrow{x}_\# N^\#, x \in \{\omega, \rho, \kappa, t\}\}} M^\#$
- “If an abstract 3D process is not present in $F \uparrow^\omega (N)$, then a concrete corresponding 3D process can never appear in the evolution of N ”
- To fix - termination is not guaranteed: infinite polymers can be created
- Next Refinement: retain **time** to verify quantitative timed safety properties like “a certain abstract 3D process will never be produced before 20 time units”

Future Work and Open Questions

- Systematize, compare and complete the existing different approaches to spatial features:
 - Deterministic Time vs Stochastic Time (through abstraction?)
 - Probabilistic Behaviours, where probability may depend on the geometrical features
- Make a clear hierarchy of spatial calculi, relating them by abstractions and/or encodings
- Investigate relationships of spatial calculi with hybrid systems for which reachability is semi-decidable
- Can the movement function be less abstract, e.g. expressed by an appropriate transition relation?
- Is there any suitable equivalence for spatial calculi at different levels of abstractions?
- Can other verification techniques be applied for (fully detailed) spatial calculi?

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