

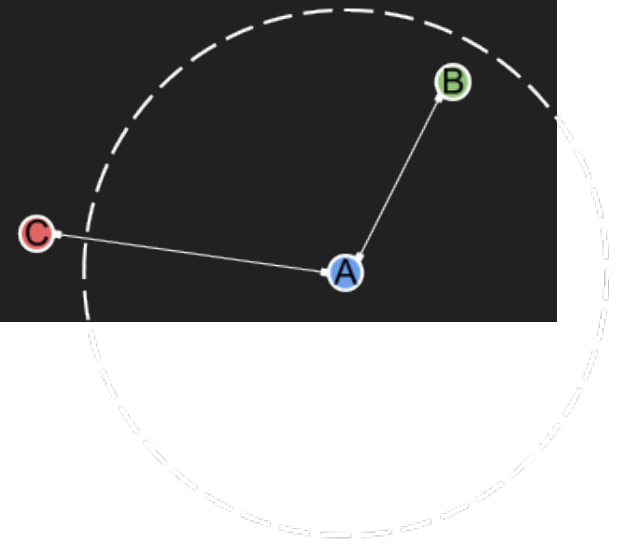
# **Privacy-preserving Location Proximity**

**Per Hallgren, Chalmers Univ. Gothenburg**

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**Andrei Sabelfeld, Chalmers University of Technology**

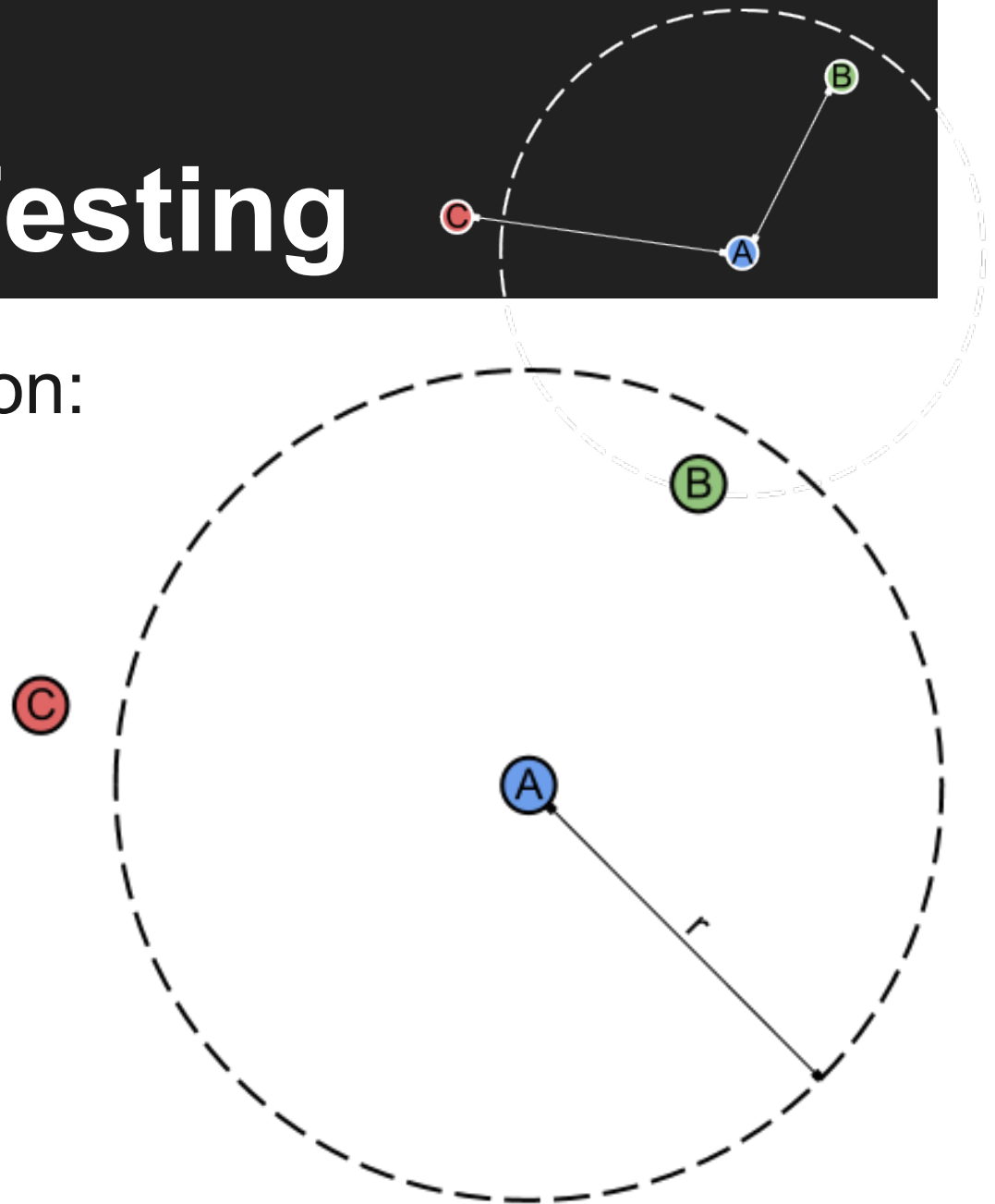
# TOC



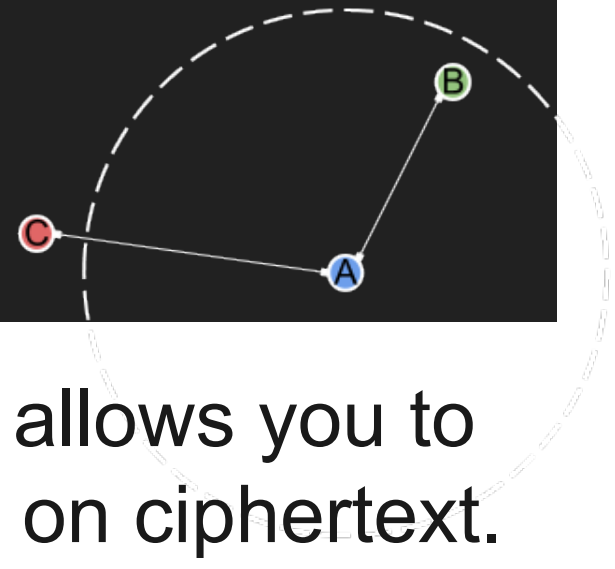
1. Background
2. Protocol
3. Theoretical Evaluation
4. Practical Evaluation

# Proximity Testing

Answers the question:  
*"Am I close?"*



# Homomorphic Encryption



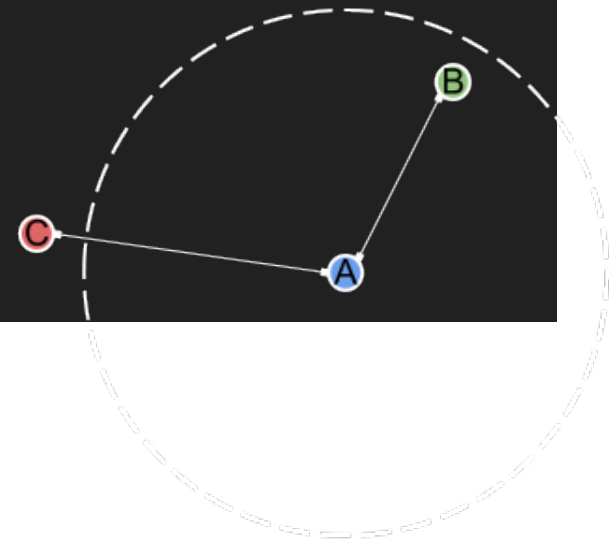
A homomorphic encryption scheme allows you to perform decipherable operations on ciphertext.

$$\text{RSA: } E(x) = x^e \bmod m$$

RSA is *multiplicatively* homomorphic

$$E(x) \times E(y) = x^e \times y^e \bmod m = (x \times y)^e \bmod m = E(x \times y)$$

# Homomorphic Encryption



Paillier:  $E(x) = g^x \bmod m$

Paillier is *additively* homomorphic

$$E(x) \times E(y) = g^x \times g^y \bmod m = g^{x+y} \bmod m = E(x+y)$$

Paillier also has this exiting property

$$E(x)^y = (g^x)^y \bmod m = g^{x \times y} \bmod m = E(x \times y)$$

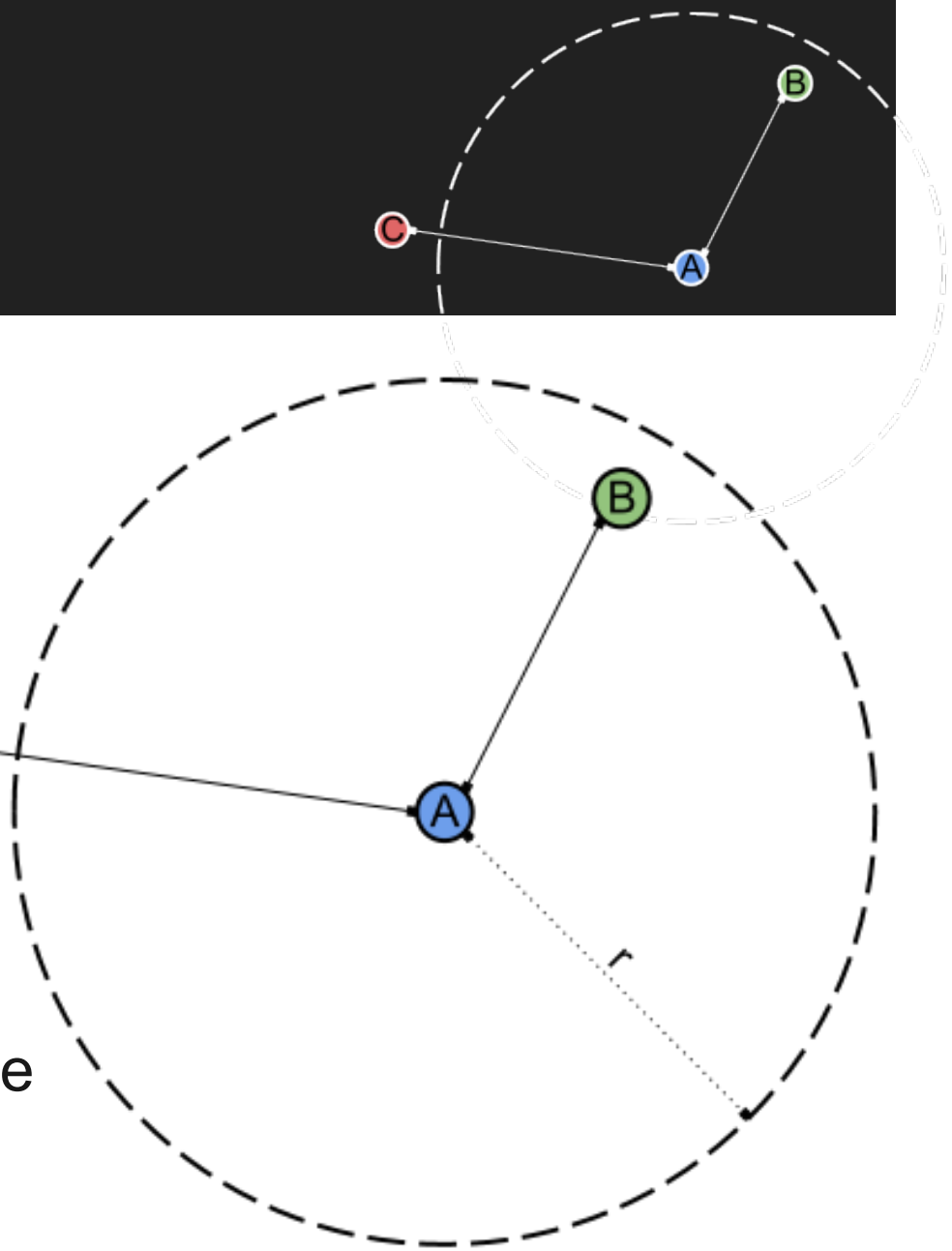
# Protocol

## Mission Statement

Answers the question:  
*"Am I close?"*

Without disclosing:

- Any information about Alice to Bob or Claire
- The position or distance of Bob and Claire to Alice



# Protocol

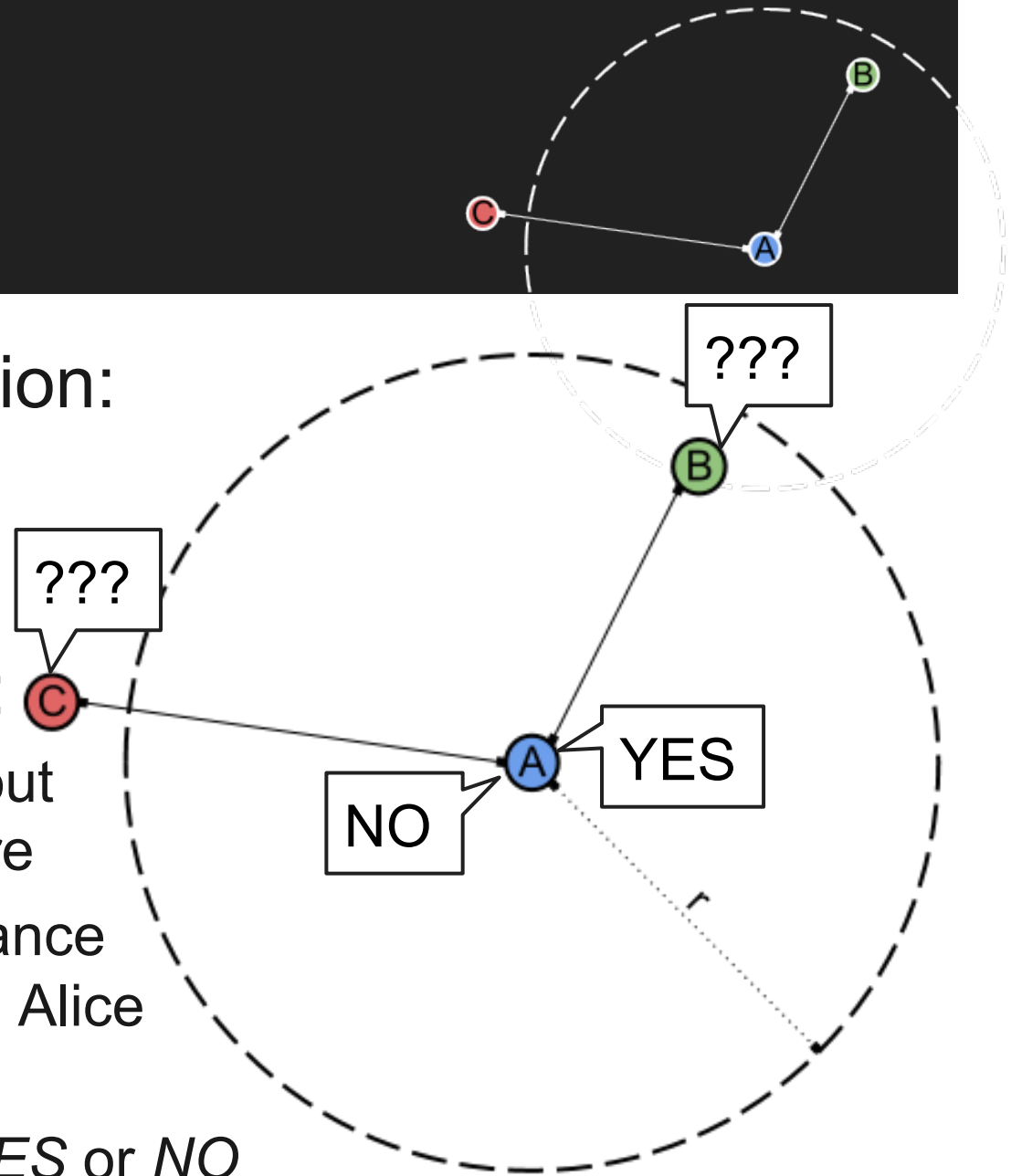
## Mission Statement

Answers the question:  
*"Am I close?"*

Without disclosing:

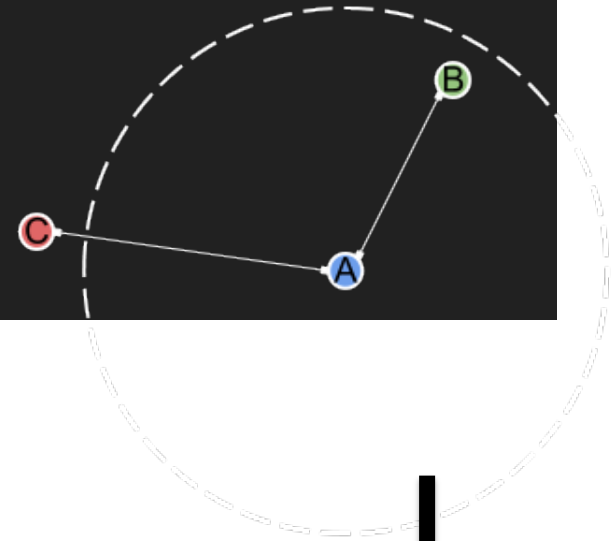
- Any information about Alice to Bob or Claire
- The position or distance of Bob and Claire to Alice

We **ONLY** say either *YES* or *NO*



# Protocol

## Outline



We **ONLY** say either *YES* or *NO*

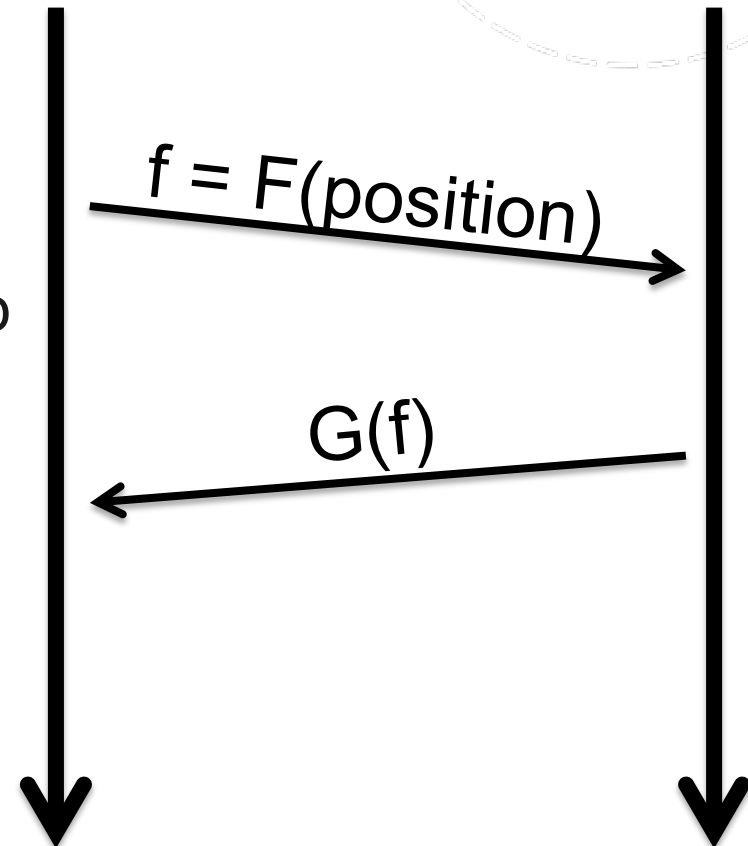
Alice

- Sends encrypted info to Bob

Bob

- Computes distance
- Sends booleanized distance

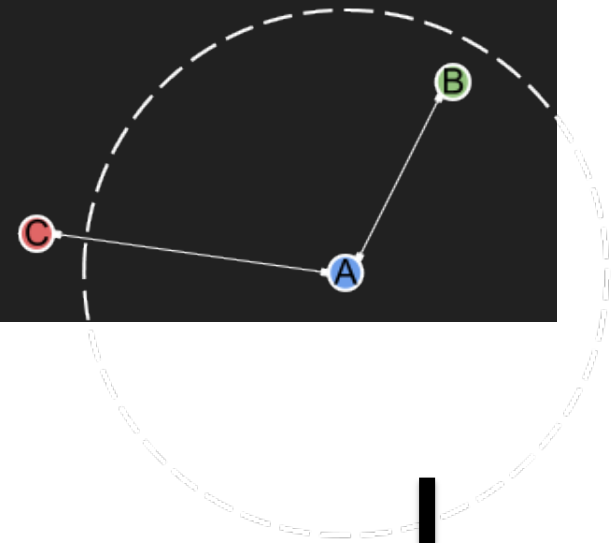
**Done!**





# Protocol

## Outline



We **ONLY** say either *YES* or *NO*

Alice:

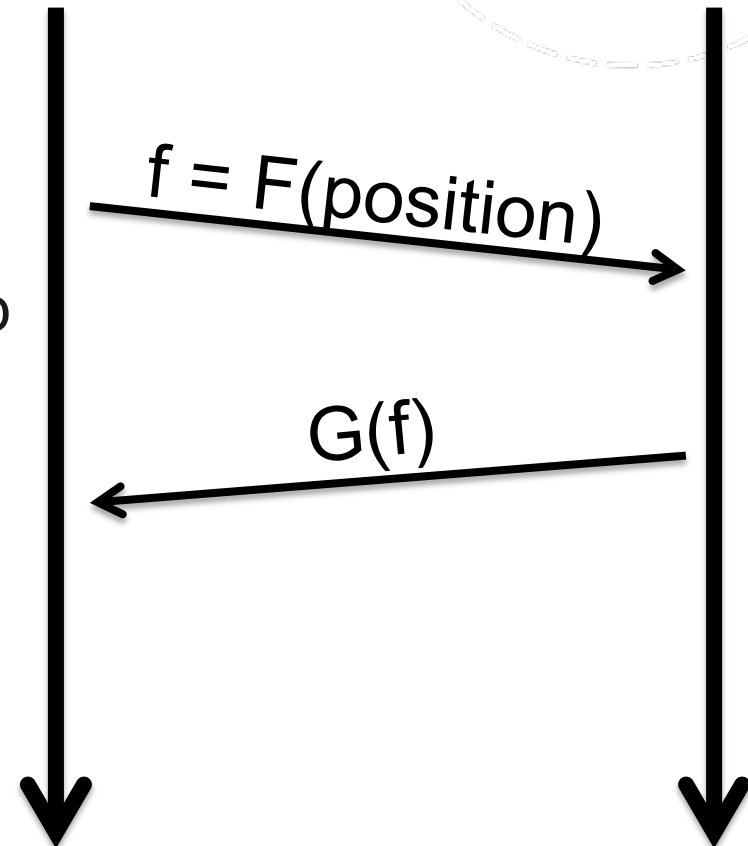
2 Sends encrypted info to Bob

Bob

1 Computes distance

3 Sends booleanized distance

**Done!**



# Protocol

## Distance Calculation

### Trivial Geometry

Distance from A to B:

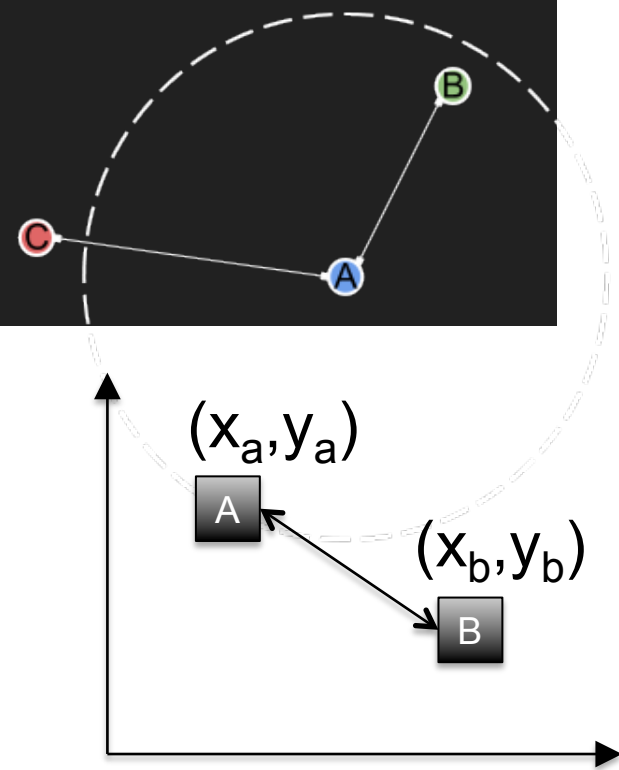
$$d = \sqrt{(x_a - x_b)^2 + (y_a - y_b)^2}$$

$$D = (x_a - x_b)^2 + (y_a - y_b)^2$$

Expand & rewrite as:

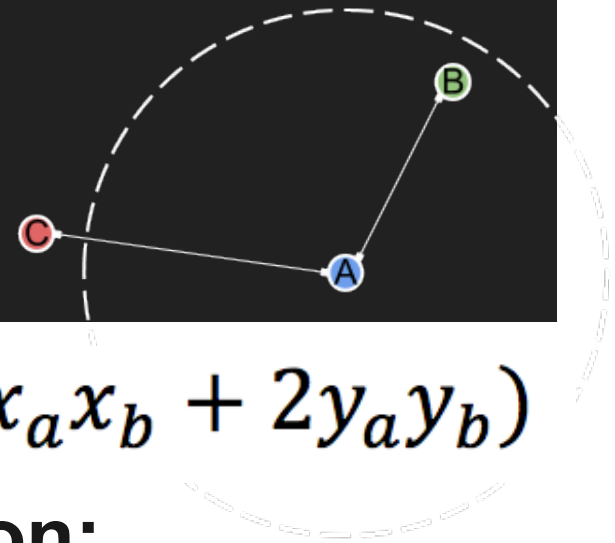
$$D = x_a^2 + x_b^2 + y_a^2 + y_b^2 - 2x_ax_b - 2y_ay_b$$

$$D = x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_ax_b + 2y_ay_b)$$



# Protocol

## Distance Calculation



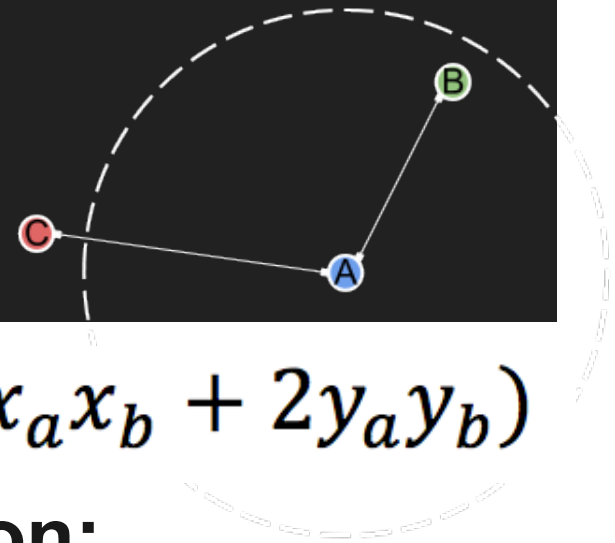
$$D = x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_ax_b + 2y_ay_b)$$

**Using Homomorphic Encryption:**

$$E(D) = E \left( x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_ax_b + 2y_ay_b) \right)$$

# Protocol

## Distance Calculation



$$D = x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_ax_b + 2y_ay_b)$$

Using Homomorphic Encryption:

$$E(D) = E \left( x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_ax_b + 2y_ay_b) \right)$$

## Recall!

Paillier is *additively* homomorphic

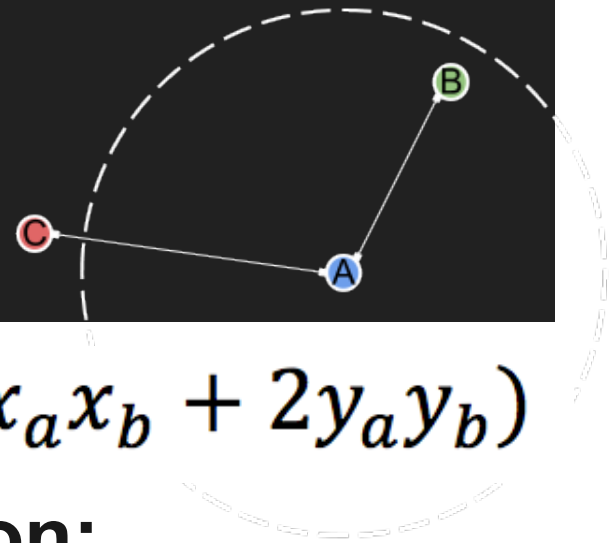
$$E(x) \times E(y) = g^x \times g^y \bmod m = E(x+y)$$

And thus:

$$E(x) / E(y) = g^x / g^y \bmod m = E(x-y)$$

# Protocol

## Distance Calculation



$$D = x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_ax_b + 2y_ay_b)$$

Using Homomorphic Encryption:

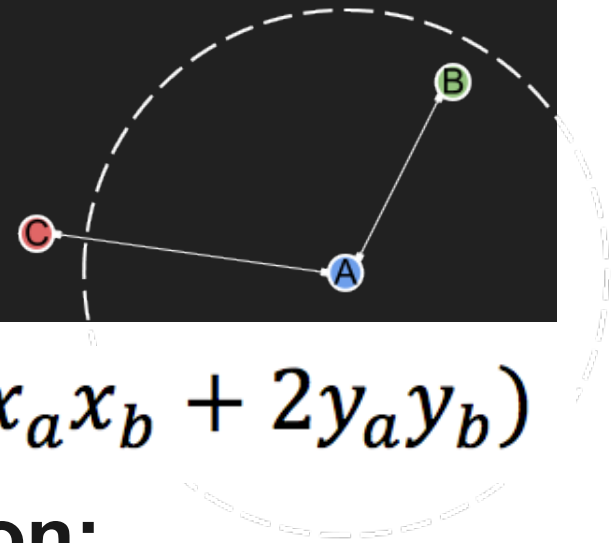
$$E(D) = E\left(x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_ax_b + 2y_ay_b)\right)$$

$$E(D) = \frac{E(x_a^2 + y_a^2 + x_b^2 + y_b^2)}{E(2x_ax_b + 2y_ay_b)}$$

$$E(D) = \frac{E(x_a^2 + y_a^2)E(x_b^2 + y_b^2)}{E(2x_ax_b)E(2y_ay_b)}$$

# Protocol

## Distance Calculation



$$D = x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_a x_b + 2y_a y_b)$$

Using Homomorphic Encryption:

$$E(D) = \frac{E(x_a^2 + y_a^2)E(x_b^2 + y_b^2)}{E(2x_a x_b)E(2y_a y_b)}$$

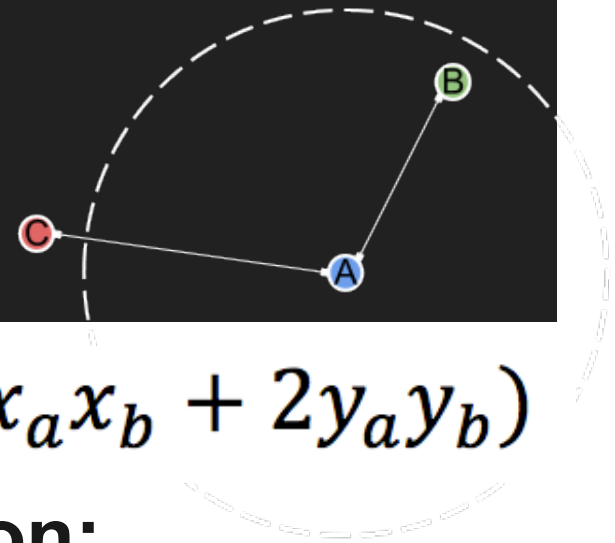
## Recall!

Raising a cipher text to a plaintext is multiplication

$$E(x)^y = (g^x)^y \bmod m = g^{x \times y} \bmod m$$

# Protocol

## Distance Calculation



$$D = x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_ax_b + 2y_ay_b)$$

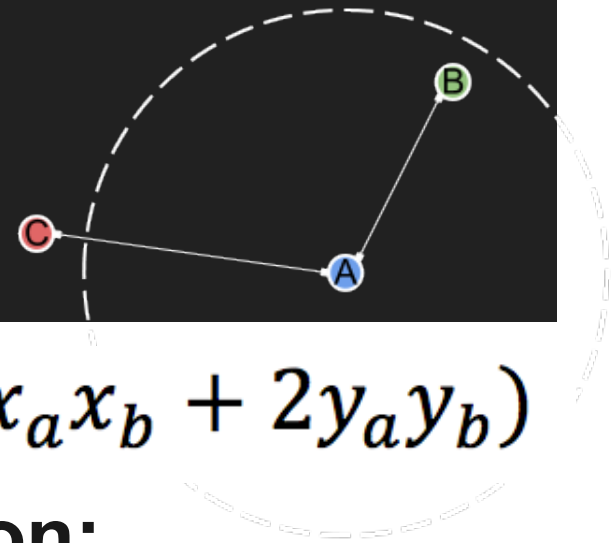
Using Homomorphic Encryption:

$$E(D) = \frac{E(x_a^2 + y_a^2)E(x_b^2 + y_b^2)}{E(2x_ax_b)E(2y_ay_b)}$$

$$E(D) = \frac{E(x_a^2 + y_a^2)E(x_b^2 + y_b^2)}{E(2x_a)^{x_b}E(2y_a)^{y_b}}$$

# Protocol

## Distance Calculation



$$D = x_a^2 + y_a^2 + x_b^2 + y_b^2 - (2x_a x_b + 2y_a y_b)$$

Using Homomorphic Encryption:

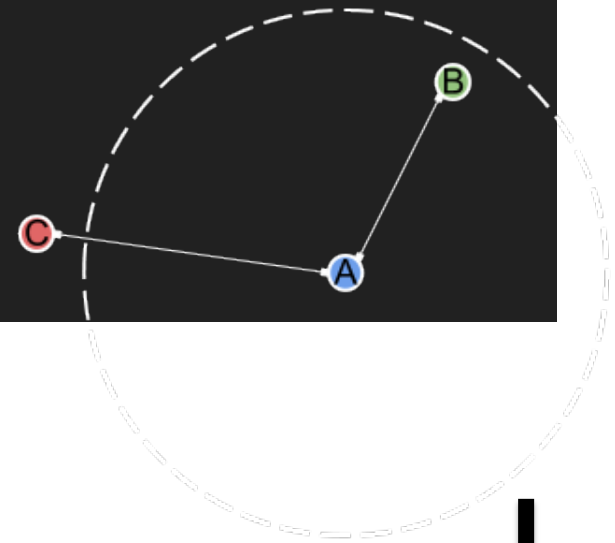
$$E(D) = \frac{E(x_a^2 + y_a^2) E(x_b^2 + y_b^2)}{E(2x_a x_b) E(2y_a y_b)}$$

$$E(D) = \frac{E(x_a^2 + y_a^2) E(x_b^2 + y_b^2)}{E(2x_a)^{x_b} E(2y_a)^{y_b}}$$



# Protocol

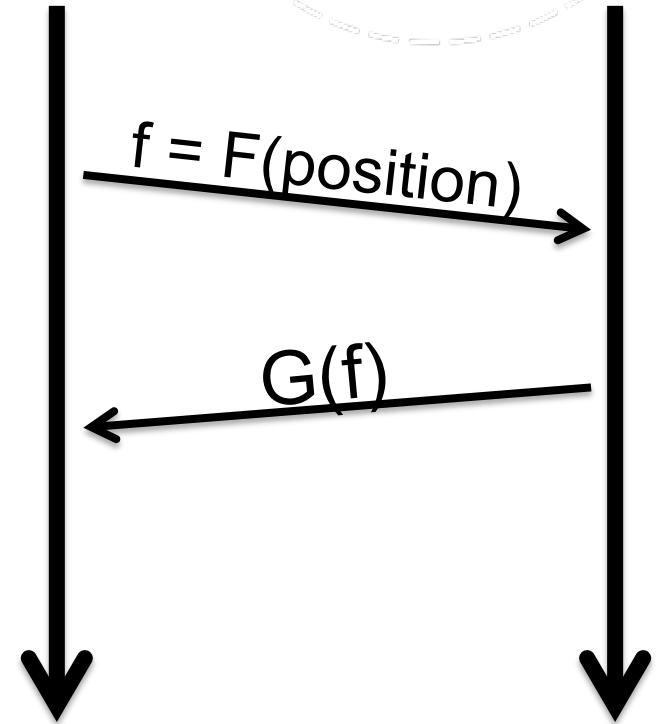
## Distance Calculation



$$E(D) = \frac{E(x_a^2 + y_a^2) E(x_b^2 + y_b^2)}{E(2x_a)^{x_b} E(2y_a)^{y_b}}$$

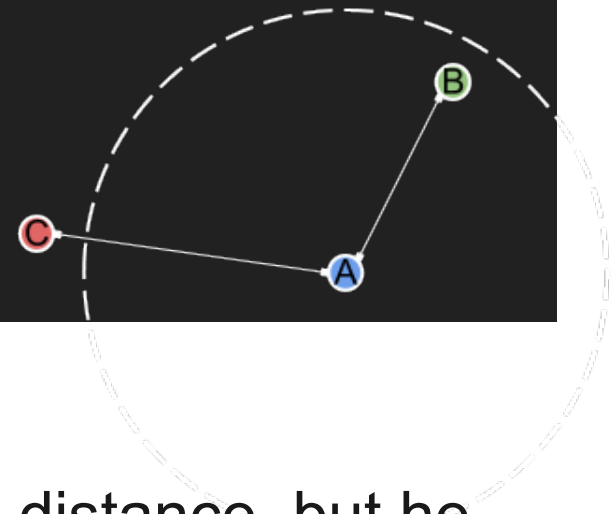
**F(position)**

$$E(x_a^2 + y_a^2) \parallel E(2x_a) \parallel E(2y_a)$$



# Protocol

## Distance Obfuscation

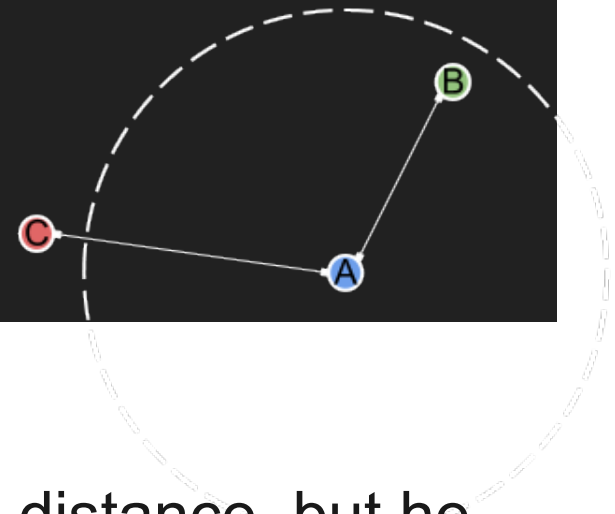


### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

# Protocol

## Distance Obfuscation



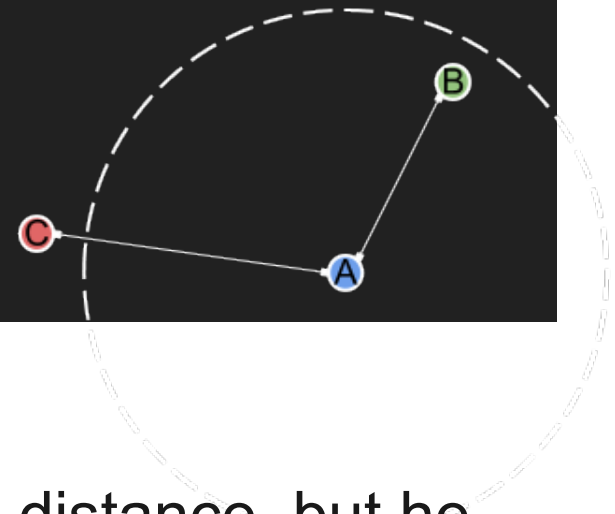
### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

- Oblivious comparison:  
 $(D-x) * \text{rand}()$

# Protocol

## Distance Obfuscation



### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

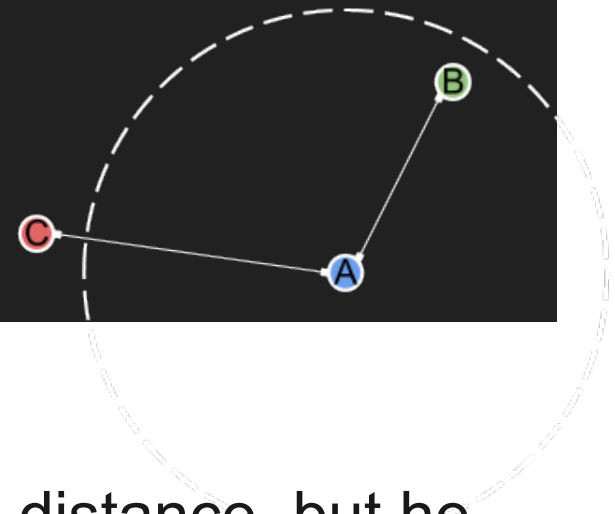
- Oblivious comparison:

$$(D-x) * \text{rand}()$$

$$\left( \frac{E_{K_A}(D)}{E_{K_A}(x)} \right)^\rho$$

# Protocol

## Distance Obfuscation



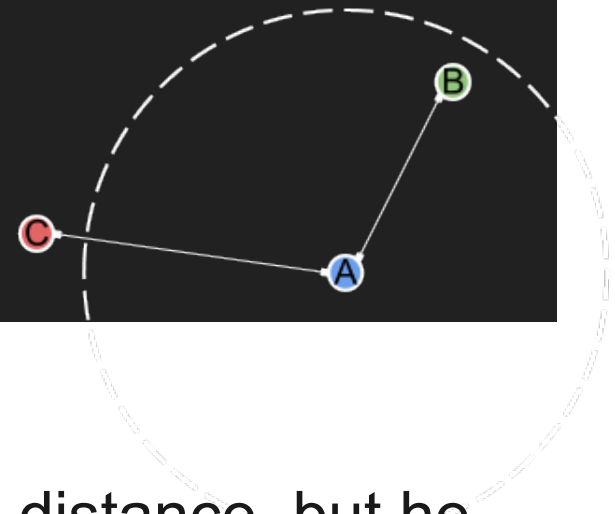
### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

- Oblivious comparison:  
 $(D-x) * \text{rand}()$
- For every  $x < r^2$ !

# Protocol

## Distance Obfuscation



### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

- Oblivious comparison:

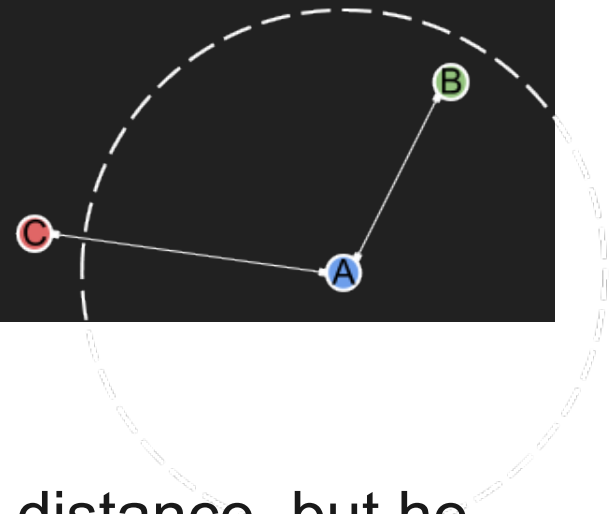
$$(D-x) * \text{rand}()$$

- For every  $x < r^2$ !

$$\left( \frac{E_{K_A}(D)}{E_{K_A}(0)} \right)^{\rho_0} :: \left( \frac{E_{K_A}(D)}{E_{K_A}(1)} \right)^{\rho_2} :: \dots :: \left( \frac{E_{K_A}(D)}{E_{K_A}(r^2)} \right)^{\rho_{r^2}}$$

# Protocol

## Distance Obfuscation



### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

- Oblivious comparison:

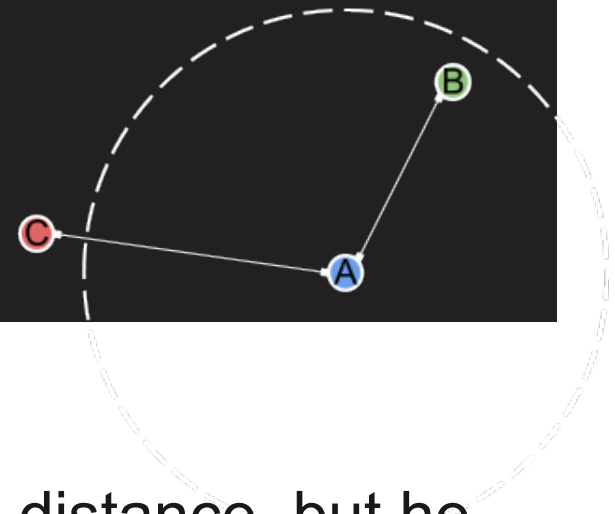
$$(D-x) * \text{rand}()$$

- For every  $x < r^2$ !
- Is this enough?

$$\left( \frac{E_{K_A}(D)}{E_{K_A}(0)} \right)^{\rho_0} :: \left( \frac{E_{K_A}(D)}{E_{K_A}(1)} \right)^{\rho_2} :: \dots :: \left( \frac{E_{K_A}(D)}{E_{K_A}(r^2)} \right)^{\rho_{r^2}}$$

# Protocol

## Distance Obfuscation



### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

- Oblivious comparison:

$$(D-x) * \text{rand}()$$

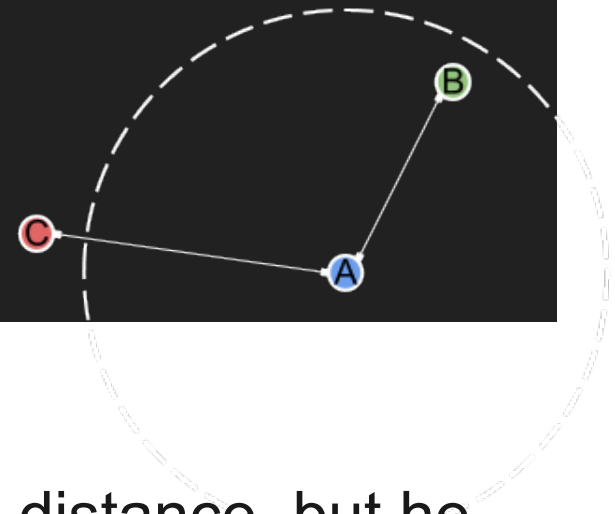
- For every  $x < r^2$ !
- Is this enough? **NO!**

$$\left( \frac{E_{K_A}(D)}{E_{K_A}(0)} \right)^{\rho_0} :: \left( \frac{E_{K_A}(D)}{E_{K_A}(1)} \right)^{\rho_2} :: \dots :: \left( \frac{E_{K_A}(D)}{E_{K_A}(r^2)} \right)^{\rho_{r^2}}$$



# Protocol

## Distance Obfuscation



### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

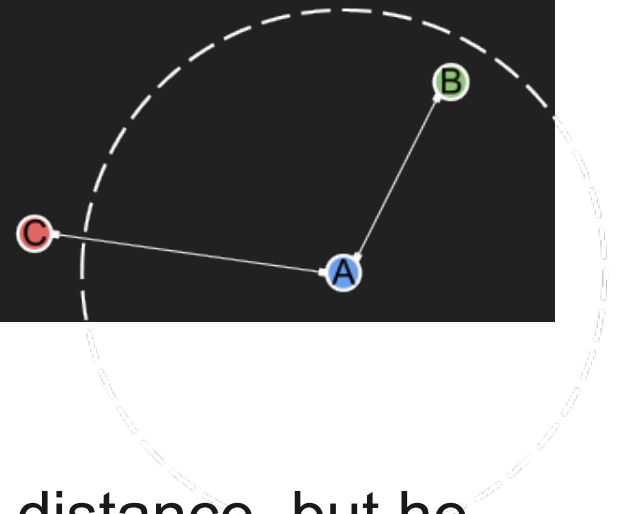
- Oblivious comparison:

$$(D-x) * \text{rand}()$$

- For every  $x < r^2$ !
- Is this enough?
- Also shuffle!

# Protocol

## Distance Obfuscation



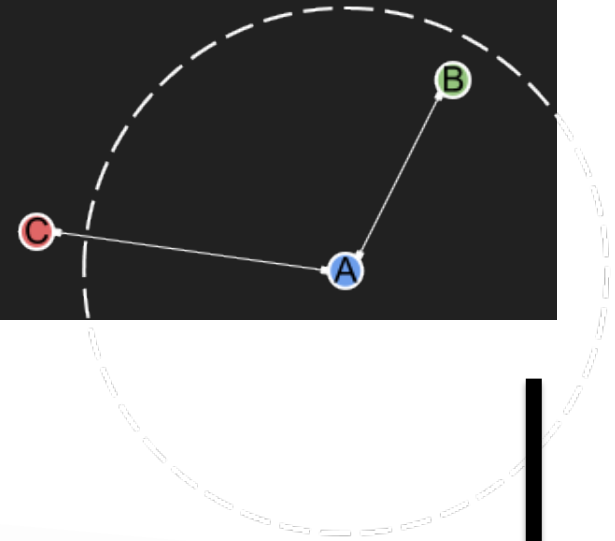
### How to obscure the distance?

Now we know how Bob can compute the distance, but he doesn't want to tell Alice what the distance is!

$$\gamma = \left( \frac{E_A(D)}{E_A(0)} \right)^{\rho_0} :: \left( \frac{E_A(D)}{E_A(1)} \right)^{\rho_2} :: \dots :: \left( \frac{E_A(D)}{E_A(r^2)} \right)^{\rho_{r^2}}$$
$$\alpha = \text{scramble}(\gamma)$$

# Protocol

## Final Result

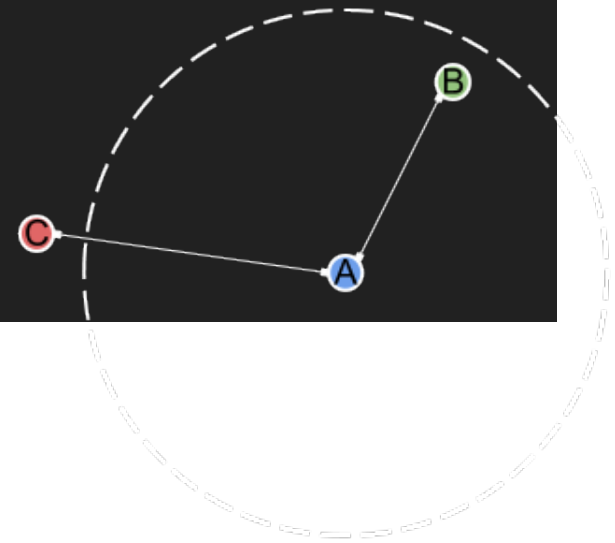


- 1: **A**  $\rightarrow$  **B**  $E_A(2x_a) :: E_A(2y_a) :: E_A(x_a^2 + y_a^2)$
- 2: **B**  $\rightarrow$  **A**  $\alpha = \text{ArrayScramble}(E_A(D))$
- 3: **A** executes **any** ( $[D_A(c) == 0 \text{ for } c \text{ in } \alpha]$ )

$E(x_a^2 + y_a^2) \parallel E(2x_a) \parallel E(2y_a)$

$\text{ArrayScramble}(E_{K_A}(D))$

# Theoretical Evaluation



## Runtime Analysis

Paillier

Encryption:  $O(\log(n) * M(n))$

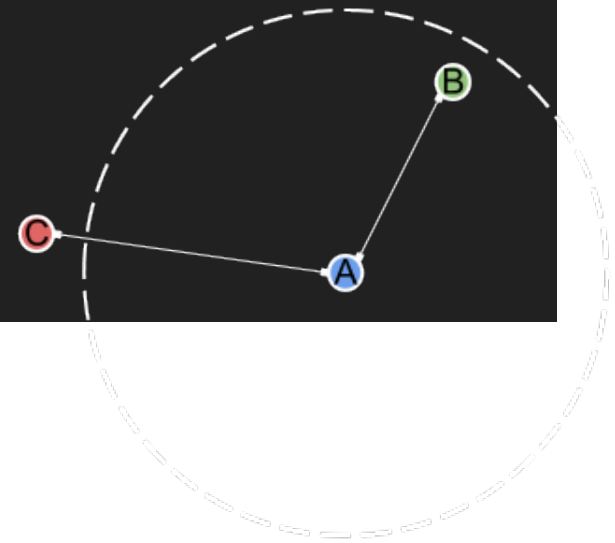
Decryption:  $O(\log(n) * M(n))$

Alice1:  $O(3\log(n) * M(n))$

Bob:  $O(r^2 * \log(n) * M(n))$

Alice2:  $O(r^2 * \log(n) * M(n))$

# Theoretical Evaluation

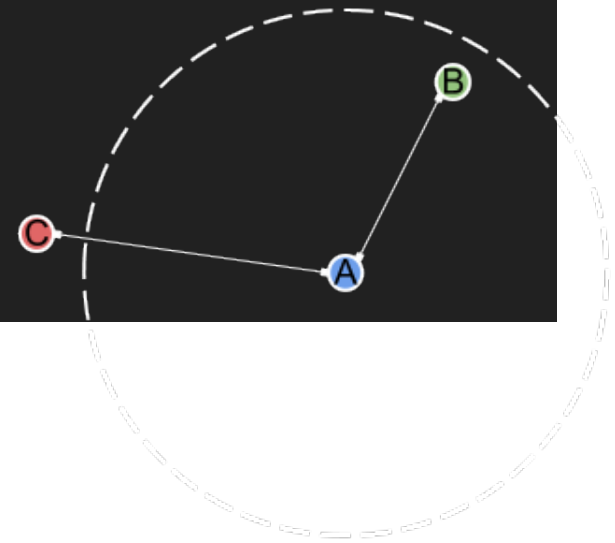


## Size Analysis

Paillier ciphertext:  $O(\log(n))$

Size of response from Bob:  $O(r^2 * \log(n))$

# Practical Evaluation



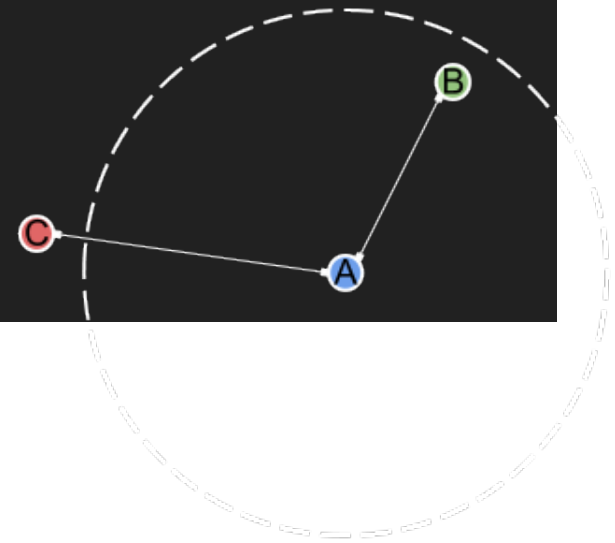
## **Proof of concept**

Small server-client application

Server relays messages to appropriate clients

All clients are interested in each other

# Practical Evaluation

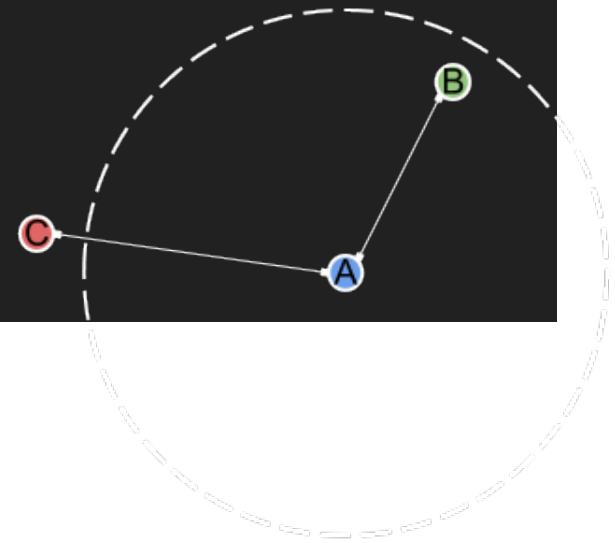


## Benchmarks

80 bit key

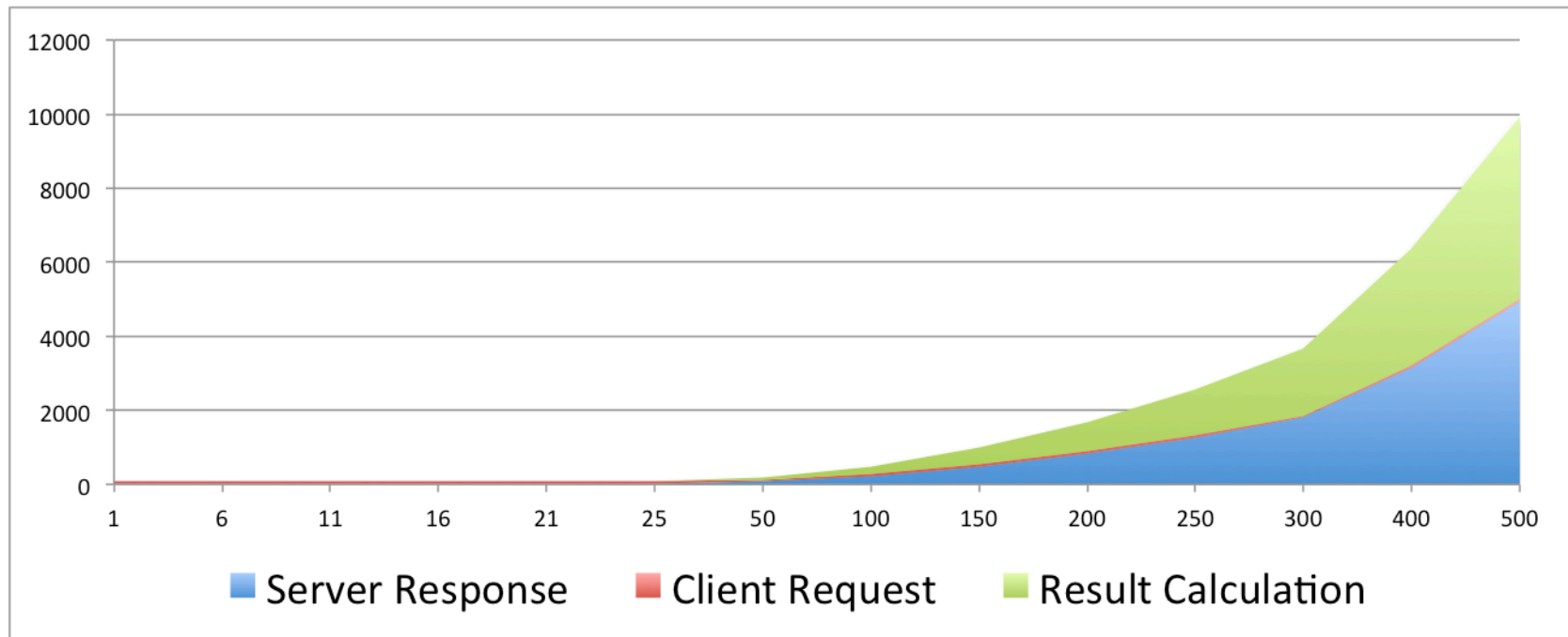


# Practical Evaluation



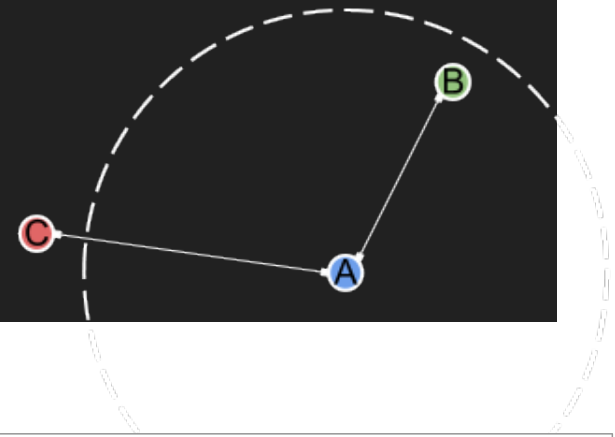
## Benchmarks

1024 bit key



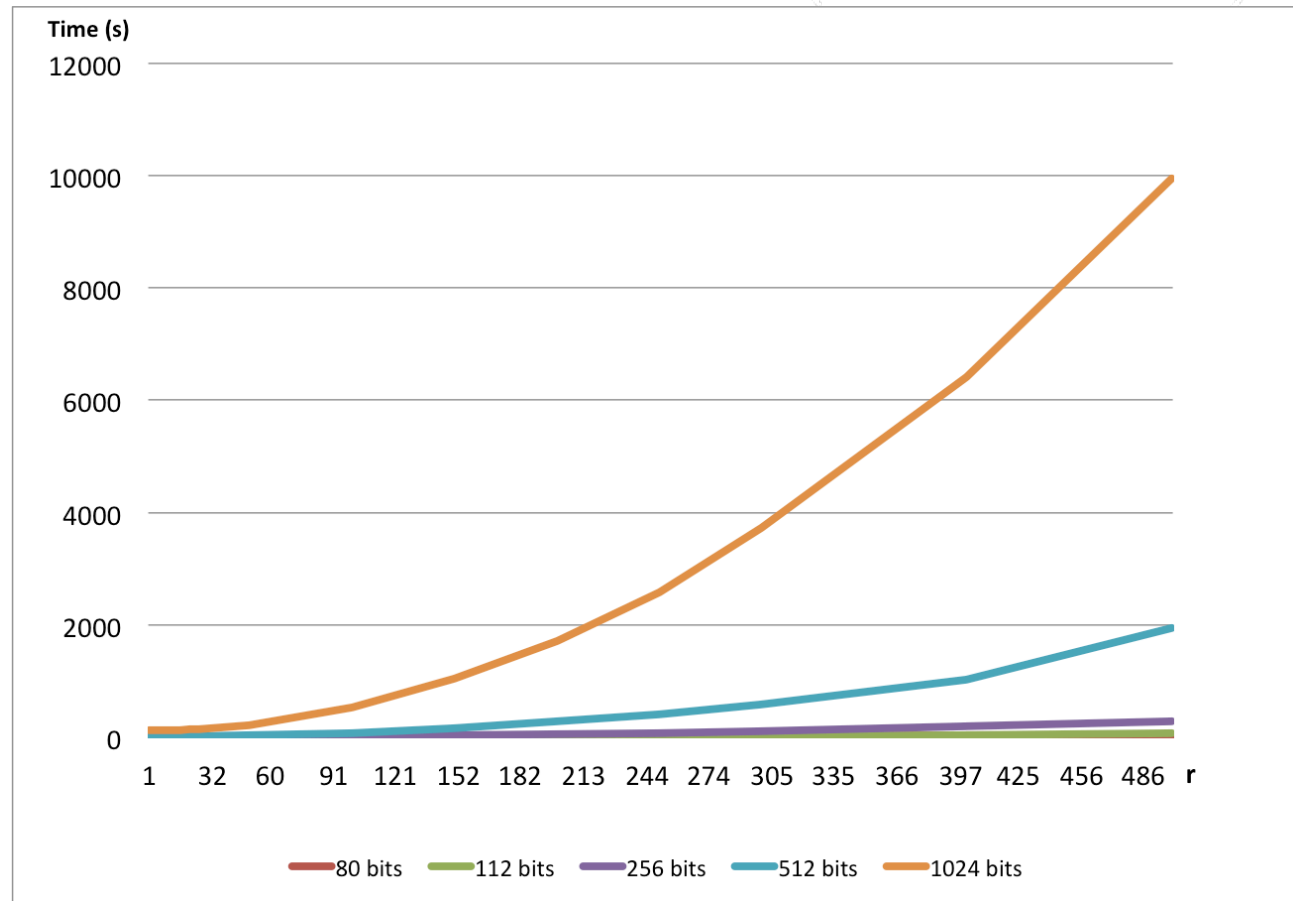


# Practical Evaluation

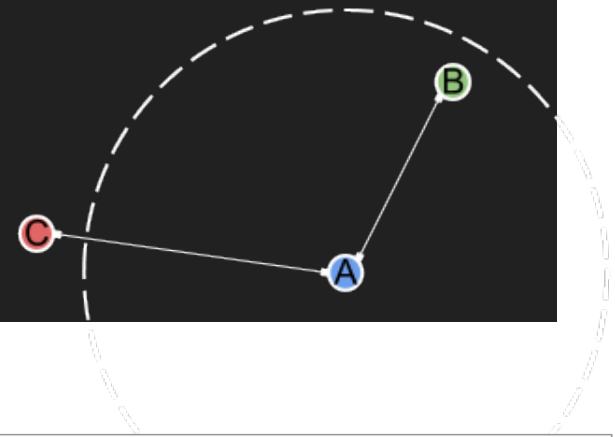


## Benchmarks

Keysize comparison



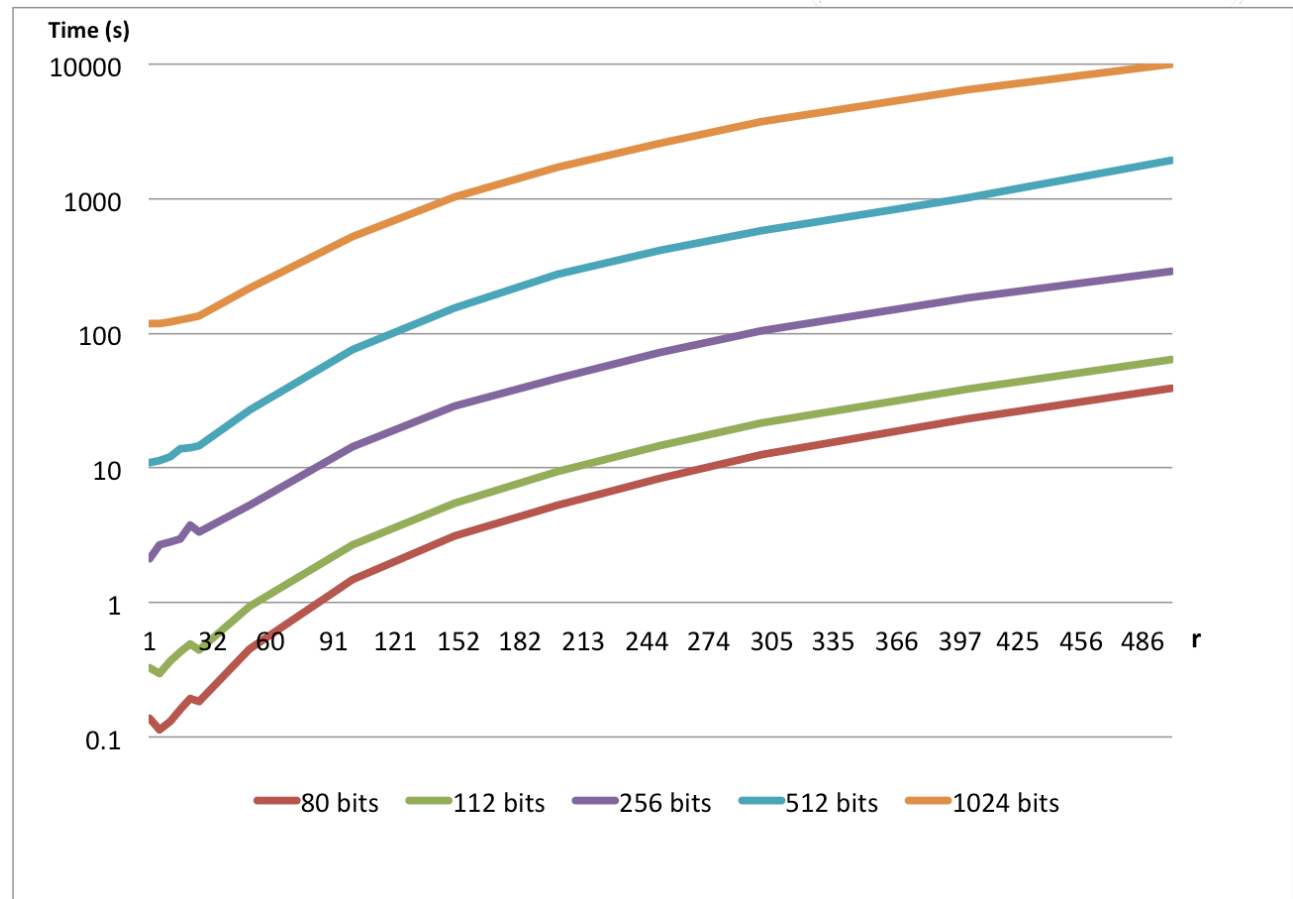
# Practical Evaluation



## Benchmarks

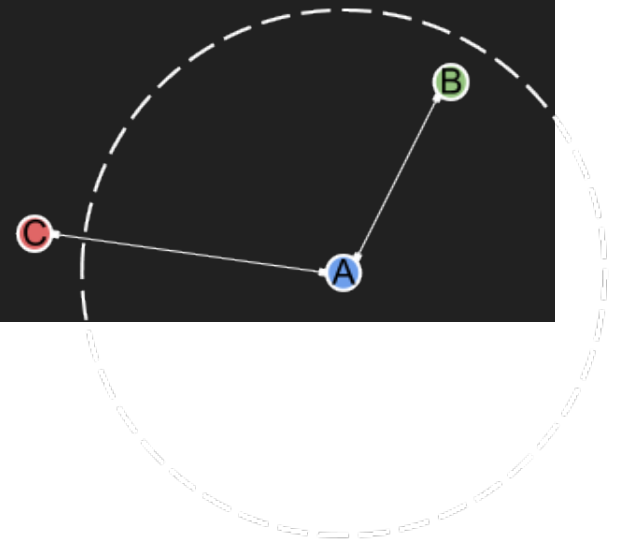
Keysize comparison

*Log scale*



# Thank You!

Questions?



# Thank You!

